

Article

Efficient Associate Rules Mining Based on Topology for Items of Transactional Data

Bo Li ^{1,*}, Zheng Pei ², Chao Zhang ^{3,*} and Fei Hao ⁴¹ School of Computer and Software Engineering, Xihua University, Chengdu 610039, China² School of Science, Xihua University, Chengdu 610039, China; zhengpei@263.com³ Intelligent Policing Key Laboratory of Sichuan Province, Sichuan Police College, Luzhou 646000, China⁴ School of Computer Science, Shaanxi Normal University, Xi'an 710119, China; fhao@snnu.edu.cn

* Correspondence: libo@mail.xhu.edu.cn (B.L.); galoiszhang1@scpolicec.edu.cn (C.Z.)

Abstract: A challenge in association rules' mining is effectively reducing the time and space complexity in association rules mining with predefined minimum support and confidence thresholds from huge transaction databases. In this paper, we propose an efficient method based on the topology space of the itemset for mining associate rules from transaction databases. To do so, we deduce a binary relation on itemset, and construct a topology space of itemset based on the binary relation and the quotient lattice of the topology according to transactions of itemsets. Furthermore, we prove that all closed itemsets are included in the quotient lattice of the topology, and generators or minimal generators of every closed itemset can be easily obtained from an element of the quotient lattice. Formally, the topology on itemset represents more general associative relationship among items of transaction databases, the quotient lattice of the topology displays the hierarchical structures on all itemsets, and provide us a method to approximate any template of the itemset. Accordingly, we provide efficient algorithms to generate Min-Max association rules or reduce generalized association rules based on the lower approximation and the upper approximation of a template, respectively. The experiment results demonstrate that the proposed method is an alternative and efficient method to generate or reduce association rules from transaction databases.

Keywords: knowledge discovery in database (KDD); frequent itemsets; closed itemsets; association rules; the topology for itemsets

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1. Introduction

In knowledge discovery in database (KDD), association rules mining (ARM) from transaction databases proposed in [1,2] have received considerable attention and wide applications, such as medical diagnosis [3,4], marketing planning [5,6], etc. ARM can be formally explained as follows: Let $\mathcal{A} = (U, A)$ be a transaction database, where U a non-empty finite set of transactions, A a non-empty finite set of items, each transaction $u_i \in U$ decides a mapping $u_i : A \rightarrow \{0, 1\}$, i.e., for any item $a_j \in A$, the transaction u_i has (or has not) the item $a_j \in A$ if $u_i(a_j) = 1$ (or $u_i(a_j) = 0$), each subset of A is called as an itemset. An association rule describing the co-occurrence relation among items is an implication in the form $A_1 \rightarrow A_2$, where itemsets A_1 and A_2 of A such that $A_1 \cap A_2 = \emptyset$ are called as antecedent and consequent, respectively. Theoretically, ARM from a transaction database is a NP-hard problem because itemsets A_1 and A_2 are selected from $2^{|A|}$ (powerset of A), in real world practices, the existed mining methods usually extract a large number of association rules which are difficult to handle; many kinds of association rules or mining methods have been proposed to generate association rules, although these mining methods are different, their processing is nearly the same, i.e., how

to evaluate usefulness of association rules and how to select the antecedent and consequent of association rules.

Quality measures are used to evaluate the usefulness of association rules, and various measures have been provided to discover significant and specific association rules. In a transaction database $\mathcal{A} = (U, A)$, for any $A' \subseteq A$, denote $\tau(A') = \{u_i \in U | \forall a_j \in A', u_i(a_j) = 1\}$, $\neg A'$ is the absence of A' , then the following quality measures can be defined

- Support [1,2]: $supp_{\mathcal{A}}(A_1 \cup A_2) = \frac{|\tau(A_1 \cup A_2)|}{|U|}$;
- Confidence [1,2]: $conf_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{supp_{\mathcal{A}}(A_1 \cup A_2)}{supp_{\mathcal{A}}(A_1)}$;
- Netconfidence [7]: $netc_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{supp_{\mathcal{A}}(A_1 \cup A_2) - supp_{\mathcal{A}}(A_1)supp_{\mathcal{A}}(A_2)}{supp_{\mathcal{A}}(A_1)(1 - supp_{\mathcal{A}}(A_1))}$;
- Conviction [8]: $conv_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{supp_{\mathcal{A}}(A_1)supp_{\mathcal{A}}(\neg A_2)}{supp_{\mathcal{A}}(A_1 \cup \neg A_2)}$;
- Added value [9]: $add_{\mathcal{A}}(A_1 \rightarrow A_2) = conf_{\mathcal{A}}(A_1 \rightarrow A_2) - supp_{\mathcal{A}}(A_2)$;
- Accuracy [9]: $acc_{\mathcal{A}}(A_1 \rightarrow A_2) = conf_{\mathcal{A}}(A_1 \rightarrow A_2) + conf_{\mathcal{A}}(\neg A_1 \rightarrow \neg A_2)$;
- Interestingness [10]: $inte_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{supp_{\mathcal{A}}(A_1 \cup A_2)}{supp_{\mathcal{A}}(A_1)} \times \frac{supp_{\mathcal{A}}(A_1 \cup A_2)}{supp_{\mathcal{A}}(A_2)} \times (1 - supp_{\mathcal{A}}(A_1 \cup A_2))$;
- Comprehensibility [10]: $comp_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{\log(1 + |\tau(A_2)|)}{\log(1 + |\tau(A_1 \cup A_2)|)}$;
- Lift [11]: $lift_{\mathcal{A}}(A_1 \rightarrow A_2) = \frac{supp_{\mathcal{A}}(A_1 \cup A_2)}{supp_{\mathcal{A}}(A_1)supp_{\mathcal{A}}(A_2)}$.

By combining with several quality measures to form fitness functions, ARM is transformed into optimization problems, and various optimal algorithms can be used to extract interesting or valid association rules from $\mathcal{A} = (U, A)$. In [12], schema constraints and the opportunistic confidence constraint are enforced to mine generalized association rules in the analyzed data. In [13], confidence, comprehensibility, and interestingness are considered as a multi-objective problem, and the particle swarm optimization algorithm is proposed to extract the best rules. In [10], support, comprehensibility and interestingness are considered as a multi-objective problem, and a Pareto-based genetic algorithm is used to extract some useful and interesting rules from any market-basket type database. In [14], Shafer's theory of evidence is used as two information measures for the quality evaluation of the set of frequent itemset (or frequent pattern). In [15], off-the-shelf constraint programming techniques are employed for modeling and solving a wide variety of constraint-based frequent itemsets. In [16], based on exclusive causal-leverage measure, a data mining algorithm is developed to mine the causal relation between drugs and their associated adverse drug reactions. In [17], the particle swarm optimization algorithm are provided to improve computational efficiency as well as to automatically determine suitable support and confidence threshold values of association rules. In [18], grammar-guided genetic programming models are proposed to deal with the association rule mining problem under a multi-objective perspective. In [19], a new multi-objective evolutionary model, which maximizes the comprehensibility, interestingness and performance of the objectives, is presented to mine a set of quantitative association rules with a good trade-off between interpretability and accuracy. In [20], a principal components analysis is applied to a set of measures that evaluate quantitative association rules' quality; the QARGA algorithm is provided to find out quantitative association rules from a wide variety of datasets. In [21], a new confidence degree of association rules and a discrete bi-level parametric programming are proposed to extract association rules from huge databases. In [22], redundancy analysis, sampling and multivariate statistical analysis are provided to ascertain the discovered rules and discard the non-significant rules. In [23], a systematic assessment of various numerical association rule mining methods and a meta-study of thirty numerical association rule mining algorithms are provided. The authors investigate how far the discretization techniques have been used in the numerical association rule mining methods. In [24], a multi-objective particle swarm optimization is proposed using an adaptive archive grid based on the Pareto optimal strategy for numerical association rule mining. In [25–28], these papers adopt the idea of parallelization and improves the Apriori algorithm based on the MapReduce model.

By analyzing associative relation among items, suitable itemsets can be selected to generate association rules. The most widely used itemsets are frequent itemsets and closed itemsets [15,29–32], an itemset is frequent if its support is not less than the minimum support value, an itemset is closed if and only if no proper superset of the itemset has the same support as the itemset. In [33], frequent itemsets are employed to construct taxonomy over items, and a breadth-first search is adopted to enumerate all frequent itemsets. Closed frequent itemsets can be used to uniquely determine the set of all frequent itemsets and applied to generate a condensed set of association rules [34–36]. Formally, all closed itemsets of $\mathcal{A} = (U, A)$ can be constructed as a closed itemset lattice, and CHARM-L algorithm is proposed to explicitly generate the frequent closed itemset lattice [37]. In [38], the parent-child relation of the closed itemset lattice has been exploited, and the cross-level closed itemset lattice is constructed to mine the most relevant minimal cross-level association rules. In [39], an efficient post-processing method is presented to prune redundant rules by virtue of the property of Galois connection, which inherently constrains rules with respect to objects. At present, itemsets have been widely investigated [40,41], and various itemsets and their generating algorithms for specific association rules mining have been proposed, such as expressive generalized itemsets [42], free itemsets [43], disjunction-free itemsets [44,45], non-derivable itemsets [46,47], disjunctive closed itemsets [48], etc. In fact, we notice that frequent itemsets or closed itemsets are rooted to the co-occurrence relation among items; from the mathematical point of view, topology may be a more suitable tool to express the relation among items, because a topology for the set is used to express a relation among subsets, subsets of the set are granulated as members of the topology, and the topology for the set is generated by its topology base; this means that the base for the topology is the basic relation among elements of the set. Topology-based pattern mining has been widely investigated in the rough set theory and knowledge model, such as in [49]; a nested topology on a crisp set of reference is provided to interpret a fuzzy subset; the approach provides us a unified framework for most of the fuzzy thresholding algorithms. In [50], the lower and upper approximations of the rough set theory are shown as pre-topological invariants and dual to each other, and the corresponding generalized topological closure via the upper approximation is constructed. In [51], the knowledge representation and reasoning method are proposed for identity-based spatial change; the change process is presented by a multistage graph, the binary relation model BC for identity change is defined, and qualitative reasoning for BC and combining BC with topological relations are investigated.

In this paper, we analyze topology and the topology base for items, express more general associative relations among items and extract useful association rules based on the topology and the base; all of these are inspired by our previous works on the formal concept analysis; in [52], we proposed a novel method based on the topology for attributes of a formal context to generate all formal concepts, the topology for attributes has been used to explain the associative relation among attributes of the formal context, and the formal concept lattice can be constructed by the topology for attributes based on an equivalent relation. In [53], we used the topology for attributes of multi-valued information systems to generate Min-Max association rules. In this paper, we construct the topology for the set of items of a transaction database, which can be used to explain more general associative relation among items; moreover, we prove that frequent itemsets, closed itemsets and closed frequent itemsets are included in the topology for the set of items, provide two kinds of lattice on the topology to display the hierarchical structures on itemsets, and propose efficient algorithms based on the base for the topology to generate or reduce many kinds of association rules.

The organization of this paper is as follows: In Section 2, we briefly review some basic notions used in association rules mining. In Section 3, we present the topology for the set of items and lattices on itemsets of a transactional database, in which some important properties of the topology, the quotient lattice and minimal generators of closed frequent itemsets, are provided. In Section 4, we propose several algorithms to generate or reduce

association rules, and analyze confidences of generalized association rules. In Section 5, we choose zoo, mushroom, connect-4 and chess as transactional databases to show the proposed method in generating association rules. We conclude the paper in Section 6.

2. Preliminaries

In the section, we briefly review some basic concepts used in frequent itemsets, closed itemsets or association rules mining; several concepts are also widely used in rough sets and formal concept analysis. For a uniform expression, we adopt the following set-to-set mapping in a transaction database $\mathcal{A} = (U, A)$, denote $P(U)$ and $P(A)$ as power sets of U and A , respectively, then we have

$$\begin{aligned} \tau & : P(A) \longrightarrow P(U), \tau(A_1) = \{u_i \in U | \forall a_j \in A_1 \subset A, u_i(a_j) = 1\}, \\ \gamma & : P(U) \longrightarrow P(A), \gamma(O) = \{a_j \in A | \forall u_i \in O \subset U, u_i(a_j) = 1\}. \end{aligned}$$

Based on τ and γ , we can rewrite the support of an itemset A_1 and the confidence of an association rule $A_1 \rightarrow A_2$, as follows:

$$\begin{aligned} \text{supp}(A_1) &= |\tau(A_1)| \text{ (or } \frac{|\tau(A_1)|}{|U|} \text{)}, \\ \text{conf}(A_1 \rightarrow A_2) &= \frac{|\tau(A_1 \cup A_2)|}{|\tau(A_1)|}. \end{aligned}$$

τ and γ also make us define closed itemsets and their generator concisely, i.e., $A_1 \subset A$ is a closed itemset if and only if $\gamma\tau(A_1) = A_1$, an itemset A_2 is a generator of the closed itemset A_1 if and only if $A_2 \subset A_1$ and $\gamma\tau(A_2) = A_1$, furthermore, A_2 is a minimal generator of A_1 if and only if $\nexists A_3 \subset A_2$ such that $\gamma\tau(A_3) = A_1$. Formally, τ and γ are also analyzed in rough sets and formal concept analysis [50,54–57], and the following properties of τ and γ are obviously, for any $A_1, A_2, A_3 \subset A$ and $O_1, O_2 \subset U$, (1) If $A_1 \subseteq A_2$, then $\tau(A_1) \supseteq \tau(A_2)$; (2) If $O_1 \subseteq O_2$, then $\gamma(O_1) \supseteq \gamma(O_2)$; (3) $\tau(A_1 \cup A_2) = \tau(A_1) \cap \tau(A_2)$, $\gamma(O_1 \cup O_2) = \gamma(O_1) \cap \gamma(O_2)$; (4) $\tau(\gamma\tau(A_1)) = \tau(A_1)$, $\gamma(\tau\gamma(O_1)) = \gamma(O_1)$. In [52], we use τ and γ to induce a reflexive and transitive relation on the set of attributes from a formal context, the topology for attributes are constructed by using the relation, then the formal concepts are analyzed in the the topology for attributes. Here, we provide the same results in a transaction database.

Definition 1. For any transactional database $\mathcal{A} = (U, A)$, τ and γ decide a point-to-set mapping from A to $P(A)$, i.e., for any $a_i \in A$, $C : A \longrightarrow P(A)$ is

$$C(a_i) = \{a_j \in A | a_j \in \gamma\tau(\{a_i\})\}.$$

Intuitively, the point-to-set mapping C represents a co-occurrence relation among items, i.e., the following binary relation on A can be decided by the mapping, C . For any $a_i, a_j \in A$,

$$R_A(a_i, a_j) = \begin{cases} 1, & \text{if } a_j \in C(a_i), \\ 0, & \text{if } a_j \notin C(a_i). \end{cases} \tag{1}$$

In [52], we have proved that the binary relation R_A on A is reflexive and transitive, and (A, R_A) is an approximation space. For any subset A_1 of A , denote $\underline{R}_A(A_1) = \{a_i \in A | \forall a_j, (R_A(a_i, a_j) = 1) \rightarrow (a_j \in A_1)\}$ and $\overline{R}_A(A_1) = \{a_i \in A | \exists a_j \in A_1 \wedge R_A(a_i, a_j) = 1\}$, then $\overline{R}_A(A_1)$ and $\underline{R}_A(A_1)$ are generalized upper approximation and lower approximation of A_1 , especially if R_A is an equivalent relation; $\overline{R}_A(A_1)$ and $\underline{R}_A(A_1)$ are Pawlak’s upper approximation and lower approximation [55]. More important, we have the following theorem.

Theorem 1 ([52]). For any transactional database $\mathcal{A} = (U, A)$.

1. $T(R_A) = \{R_A(A_1) \mid A_1 \subseteq A\}$ is a topology for A ;
2. For any $a_i \in A$, $R_A(C(a_i)) = C(a_i)$;
3. $\mathcal{B}_A = \{C(a_i) \mid a_i \in A\}$ is a base for the topology $T(R_A)$.

The theorem means that the topology for A has a simple expression, i.e., $T(R_A) = \{\cup_{C(a_i) \in \mathcal{B}'_A} C(a_i) \mid \forall \mathcal{B}'_A \subseteq \mathcal{B}_A\}$, we use the following example to explain the above mentioned concepts and results; the example was initially used to analyze the concept “to-be-a-fruit” [58].

Example 1. Let a fruit data be $\mathcal{A} = (U, A)$, where $U = \{a \text{ chestnut } (u_1), \text{ an olive } (u_2), \text{ a pepper } (u_3), \text{ a strawberry } (u_4), \text{ an orange } (u_5), \text{ a tomato } (u_6)\}$ and $A = \{\text{to-grow-on-trees } (a_1), \text{ to-be-sweet } (a_2), \text{ to-be-raw-edible } (a_3), \text{ to-yeild-juice } (a_4), \text{ to-have-a-skin } (a_5)\}$ be used to understand the concept “to-be-a-fruit” (shown in Table 1). According to Definition 1, we have $C(a_1) = \{a_1\}$, $C(a_2) = \{a_2, a_3\}$, $C(a_3) = \{a_3\}$, $C(a_4) = \{a_3, a_4, a_5\}$ and $C(a_5) = \{a_5\}$, e.g., for $C(a_2) = \{a_2, a_3\}$, it means that a_2 and a_3 are co-occurrence in Table 1, i.e., a fruit is “to-be-sweet”, it must be “to-be-raw-edible”. The binary relation R_A on $A = \{a_1, a_2, a_3, a_4, a_5\}$ is shown in Table 2.

Table 1. A transactional database $\mathcal{A} = (U, A)$.

$U \setminus A$	a_1	a_2	a_3	a_4	a_5
u_1	1	0	0	0	1
u_2	1	0	1	0	0
u_3	0	0	1	0	1
u_4	0	1	1	0	0
u_5	1	1	1	1	1
u_6	0	0	1	1	1

Table 2. The binary relation R_A on A .

R_A	a_1	a_2	a_3	a_4	a_5
a_1	1	0	0	0	0
a_2	0	1	1	0	0
a_3	0	0	1	0	0
a_4	0	0	1	1	1
a_5	0	0	0	0	1

Based on Theorem 1 and the mapping C , we have $T(R_A) = \{\emptyset, \{a_1\}, \{a_3\}, \{a_5\}, \{a_1, a_3\}, \{a_1, a_5\}, \{a_2, a_3\}, \{a_3, a_5\}, \{a_1, a_2, a_3\}, \{a_1, a_3, a_5\}, \{a_2, a_3, a_5\}, \{a_3, a_4, a_5\}, \{a_1, a_3, a_4, a_5\}, \{a_1, a_2, a_3, a_5\}, \{a_2, a_3, a_4, a_5\}, \{a_1, a_2, a_3, a_4, a_5\}\}$, e.g., the member $\{a_2, a_3, a_5\}$ is obtained by

$$\{a_2, a_3, a_5\} = C(a_2) \cup C(a_5).$$

According to Table 2, we can find that R_A on A is not symmetric generally, e.g., $R_A(a_2, a_3) = 1$ but $R_A(a_3, a_2) = 0$, hence, members of the topology $T(R_A)$ for the set of items are not Pawlak’s lower approximation in practice. According to Table 1, we can obtain the support of each member in the topology for A , e.g., for $\{a_2, a_3\}$, we have $\text{supp}(\{a_2, a_3\}) = |\tau(T)| = |\{u_4, u_5\}| = 2$ or $\text{supp}(\{a_2, a_3\}) = \frac{|\{u_4, u_5\}|}{|U|} = \frac{2}{3}$. If we fix minimum support $\frac{1}{3}$, then members such as $\{a_2, a_3\}$ and $\{a_1, a_3\}$ are frequent itemsets, which can be used to generate association rules, e.g., we can generate association rules “ $a_1 \rightarrow a_3$ ” and “ $a_3 \rightarrow a_1$ ”; confidences of them are

$$\begin{aligned} \text{conf}(a_1 \rightarrow a_3) &= \frac{|\tau(\{a_1, a_3\})|}{|\tau(a_1)|} = \frac{|\{u_2, u_5\}|}{|\{u_1, u_2, u_5\}|} = \frac{2}{3} \\ \text{conf}(a_3 \rightarrow a_1) &= \frac{|\tau(\{a_1, a_3\})|}{|\tau(a_3)|} = \frac{|\{u_2, u_5\}|}{|\{u_2, u_3, u_4, u_5, u_6\}|} = 0.4. \end{aligned}$$

In A of the fruit data $\mathcal{A} = (U, A)$, “ $a_1 \rightarrow a_3$ ” is “to-grow-on-trees implies to-be-raw-edible”, “ $a_3 \rightarrow a_1$ ” is “to-be-raw-edible implies to-grow-on-trees”, by considering their confidences, “ $a_1 \rightarrow a_3$ ” is more confident than “ $a_3 \rightarrow a_1$ ” due to $\frac{2}{3} > 0.4$.

From Example 1, we notice that members of the topology $T(R_A)$ for the set of items are itemsets, which is generated by the co-occurrence relation among items. In existed ARM methods, frequent itemsets or closed itemsets are always used to mine association rules, one problem is “can we use members of the topology $T(R_A)$ for the set of items to mine all useful and needful association rules”. The problem will be solved in the next section.

3. Lattice Structures on the Topology for the Set of Items

In this section, we use set inclusion to construct lattice structures on the topology $T(R_A)$ for the set of items, one is constructed on the topology $T(R_A)$ itself, another is constructed on quotient set of the topology $T(R_A)$, then we analyze minimal elements and minimal generators of closed itemsets in the lattice structures, which are useful information in ARM.

3.1. The Lattice on the Topology

Formally, the topology for the set of items is a proper subset of the power set of A —theoretically, the power set of A naturally forms a power set lattice; however, its proper subset may be not a lattice. Here, we use the set inclusion to construct lattice on the topology for the set of items and analyze hierarchical structure of members of the topology.

It is obvious that $T(R_A)$ is an poset by set inclusion, i.e., for any $T_1, T_2 \in T(R_A)$,

$$T_1 \leq T_2 \text{ if and only if } T_1 \subseteq T_2.$$

On the poset $(T(R_A), \leq)$, we define

$$T_1 \wedge T_2 = \left(\bigcup_{a_i \in A}^{C(a_i) \subseteq T_1} C(a_i) \right) \cap \left(\bigcup_{a_j \in A}^{C(a_j) \subseteq T_2} C(a_j) \right),$$

$$T_1 \vee T_2 = \left(\bigcup_{a_i \in A}^{C(a_i) \subseteq T_1} C(a_i) \right) \cup \left(\bigcup_{a_j \in A}^{C(a_j) \subseteq T_2} C(a_j) \right).$$

For any $a_k \in T_1 \wedge T_2$, we have $a_k \in \bigcup_{a_i \in A}^{C(a_i) \subseteq T_1} C(a_i)$ and $a_k \in \bigcup_{a_j \in A}^{C(a_j) \subseteq T_2} C(a_j)$; hence, there exists $C(a_i) \subseteq T_1$ and $C(a_j) \subseteq T_2$ such that $a_k \in C(a_i)$ and $a_k \in C(a_j)$, according to Theorem 1, $C(a_k) \subseteq C(a_i)$ and $C(a_k) \subseteq C(a_j)$, we have

$$C(a_k) \subseteq \left(\bigcup_{a_i \in A}^{C(a_i) \subseteq T_1} C(a_i) \right) \cap \left(\bigcup_{a_j \in A}^{C(a_j) \subseteq T_2} C(a_j) \right)$$

$$= \bigcup_{a_i \in A}^{C(a_i) \subseteq T_1} \bigcup_{a_j \in A}^{C(a_j) \subseteq T_2} (C(a_i) \cap C(a_j)) = T_1 \wedge T_2,$$

this means that we rewrite $T_1 \wedge T_2 = \bigcup_{a_k \in T_1 \cap T_2} C(a_k) \in T(R_A)$. Similarly, we rewrite $T_1 \vee T_2 = \bigcup_{a_k \in T_1 \cup T_2} C(a_k) \in T(R_A)$, this means that $(T(R_A), \wedge, \vee)$ is a lattice.

Furthermore, for any subset $\mathcal{T} \subseteq T(R_A)$, we denote

$$\begin{aligned} \bigwedge \mathcal{T} &= \bigwedge_{T_i \in \mathcal{T}} T_i = \bigcap_{T_i \in \mathcal{T}} \left(\bigcup_{a_i \in A}^{C(a_i) \subseteq T_i} C(a_i) \right), \\ \bigvee \mathcal{T} &= \bigvee_{T_i \in \mathcal{T}} T_i = \bigcup_{T_i \in \mathcal{T}} \left(\bigcup_{a_i \in A}^{C(a_i) \subseteq T_i} C(a_i) \right), \end{aligned}$$

then, one can easily prove that $\bigwedge \mathcal{T} = \bigcup_{a_k \in \bigcap_{T_i \in \mathcal{T}} T_i} C(a_k) \in T(R_A)$ and $\bigvee \mathcal{T} = \bigcup_{a_k \in \bigcup_{T_i \in \mathcal{T}} T_i} C(a_k) \in T(R_A)$, and $(T(R_A), \wedge, \vee)$ is also a complete lattice, e.g., in Example 1, we have $\{a_1\} \wedge \{a_3\} = \emptyset$, $\{a_1, a_2, a_3\} \wedge \{a_1, a_3, a_5\} = C(a_1) \cup C(a_3) = \{a_1, a_3\}$ due to $\{a_1, a_2, a_3\} \cap \{a_1, a_3, a_5\} = \{a_1, a_3\}$, and $\{a_2, a_3, a_5\} \vee \{a_3, a_4, a_5\} = C(a_2) \cup C(a_3) \cup C(a_4) \cup C(a_5) = \{a_2, a_3, a_4, a_5\}$ due to $\{a_2, a_3, a_5\} \cup \{a_3, a_4, a_5\} = \{a_2, a_3, a_4, a_5\}$. The complete lattice $(T(R_A), \wedge, \vee)$ of Example 1 is shown in Figure 1.

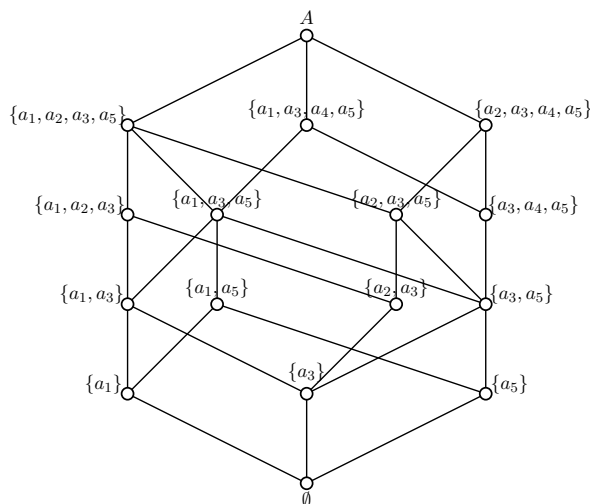


Figure 1. The lattice of the topology $T(R_A)$ of Example 1.

From up to down, the complete lattice $(T(R_A), \wedge, \vee)$ provides us hierarchical information among members of the topology, intuitively, hierarchical information means that any member of $T(R_A)$ is contained in its upper hierarchical members, and there are common transactions in the member of $T(R_A)$ and its upper hierarchical members; such information can be used to generate association rules, i.e., suppose that T_2 is an upper hierarchical member of T_1 , then, we have the following association rule

$$\psi = T_1 \rightarrow (T_2 - T_1),$$

in which, $supp(\psi) = |\tau(T_2)|$ and $conf(\psi) = \frac{|\tau(T_2)|}{|\tau(T_1)|}$, e.g., in Figure 1, itemset $\{a_3, a_5\}$ is an up hierarchical member of itemset $\{a_5\}$, according to Table 1, we have $a_5 \rightarrow a_3$ with $supp(a_5 \rightarrow a_3) = |\tau(\{a_3, a_5\})| = 3$ and $conf(a_5 \rightarrow a_3) = \frac{|\tau(\{a_3, a_5\})|}{|\tau(\{a_5\})|} = 0.75$. $a_1 \wedge a_3 \rightarrow a_5$ with $supp(a_1 \wedge a_3 \rightarrow a_5) = |\tau(\{a_1, a_3, a_5\})| = 1$ and $conf(a_1 \wedge a_3 \rightarrow a_5) = \frac{|\tau(\{a_1, a_3, a_5\})|}{|\tau(\{a_1, a_3\})|} = 0.5$. In A of the fruit data $\mathcal{A} = (U, A)$, “ $a_5 \rightarrow a_3$ ” is “to-have-a-skin implies to-be-raw-edible”, “ $a_1 \wedge a_3 \rightarrow a_5$ ” is “to-grow-on-trees and to-be-raw-edible implies to-have-a-skin”; obviously, “to-have-a-skin implies to-be-raw-edible” is more confident and useful than “to-grow-on-trees and to-be-raw-edible implies to-have-a-skin” due to $0.75 > 0.5$.

3.2. The Lattice on the Quotient Set of the Topology

To fast mine association rules with high support and confidence from the topology for the set of items, we construct another lattice structure on the topology in the subsection;

then, we analyze minimal elements and minimal generators of closed itemsets in the lattice, which can help us to fast generate association rules with high support and confidence.

For any $T_1, T_2 \in T(R_A)$, we define a binary relation on the topology for the set of items, as follows

$$T_1 \sim_\tau T_2 \text{ if and only if } \tau(T_1) = \tau(T_2).$$

It is obvious that \sim_τ is an equivalent relation on the topology $T(R_A)$, and an quotient set of the topology $T(R_A)$ can be decided by the equivalent relation, i.e., $T(R_A)/\sim_\tau = \{[T] | T \in T(R_A)\}$; each equivalent class $[T]$ in $T(R_A)/\sim_\tau$ is consisted of members of $T(R_A)$ with the same support $\tau(T)$. According to the property of the topology, we have $T_1 \cup T_2$ is in $T(R_A)$ if T_1 and T_2 are in $T(R_A)$. On the other hand, according to property of the set-to-set mapping τ , for any $[T] \in T(R_A)/\sim_\tau$ and $T', T'' \in [T]$, $\tau(T' \cup T'') = \tau(T') \cap \tau(T'') = \tau(T)$, we have $T' \cup T'' \in [T]$, i.e., each equivalent class $[T]$ in $T(R_A)/\sim_\tau$ is closed for the \cup operation, hence, we have the following theorem.

Theorem 2. For each equivalent class $[T]$ in $T(R_A)/\sim_\tau$, $[T]$ is a union semi-lattice.

The maximum element of each equivalent class $[T]$ in $T(R_A)/\sim_\tau$ is important. Formally, we denote the maximum element of $[T]$ as $\cup[T]$, then we confirm that $\cup[T]$ is a closed itemset and each member T of $[T]$ is a generator of $\cup[T]$, in fact, suppose that T is a generator of a closed itemset A' , according to the set-to-set mapping τ and Definition 1; we have

$$\tau(A') = \bigcap_{a_i \in A'} \tau(a_i) = \bigcap_{a_i \in A'} \tau(C(a_i)) = \tau(\bigcup_{a_i \in A'} C(a_i)) = \tau(T),$$

this means that A' is a member of the topology for the set of items and $A' \in [T]$, hence, $A' = \cup[T]$.

Corollary 1. For any transactional database $\mathcal{A} = (U, A)$ and each equivalent class $[T]$ in $T(R_A)/\sim_\tau$, (1) all closed itemsets of A are in $T(R_A)$; (2) $\cup[T]$ is a closed itemset; (3) $\forall T' \in [T]$ and $T' \neq \cup[T]$, T' is a generator of $\cup[T]$.

Because $T(R_A)$ is a proper subset of powerset of A and the existed methods search closed itemsets in powerset of A , the corollary means that we can reduce the searching range of closed itemsets in the topology for the set of items. On the other hand, frequent itemsets own the downward closed property, i.e., any subset of a frequent itemset is still frequent; hence, we can use closed frequent itemsets to obtain its all frequent itemsets, any subset of a closed frequent itemset is a frequent itemset. The following corollaries help us to obtain minimal generators of closed itemsets, which is useful information in ARM.

In each union semi-lattice $[T]$, we denote minimal members of $[T]$ as $min[T] = \{T' \in [T] | \nexists T'' \in [T] \wedge T'' \subseteq T'\}$.

Corollary 2. For any $T' \in min[T]$, if there exists no $P \subset T'$ such that $\tau(P) = \tau(T')$, then T' is a minimal generator of $T'' \in [T]$ such that $T'' \notin min[T]$ and $T' \subset T''$.

In Example 1, we can easily check $\{[a_1, a_2, a_3]\} = \{\{a_1, a_2, a_3\}, \{a_1, a_3, a_5\}, \{a_2, a_3, a_5\}, \{a_1, a_3, a_4, a_5\}, \{a_1, a_2, a_3, a_5\}, \{a_2, a_3, a_4, a_5\}, \{a_1, a_2, a_3, a_4, a_5\}\}$, in which, $min\{[a_1, a_2, a_3]\} = \{\{a_1, a_2, a_3\}, \{a_1, a_3, a_5\}, \{a_2, a_3, a_5\}\}$, due to for any $P \subset \{a_1, a_3, a_5\}$, $\tau(P) \neq \tau(\{a_1, a_3, a_5\}) = \{u_5\}$, $\{a_1, a_3, a_5\}$ is a minimal generator of $\{a_1, a_3, a_4, a_5\}$, $\{a_1, a_2, a_3, a_5\}$ and $\{a_1, a_2, a_3, a_4, a_5\}$.

For $T' \in min[T]$, if there exists $P \subset T'$ such that $\tau(P) = \tau(T')$, we denote $\mathcal{P}_\tau(T') = \{P \subset T' | T' \in min[T], \tau(P) = \tau(T')\}$, $\mathcal{P}_\tau(T')$ is an poset by set inclusion, i.e., $\forall P_1, P_2 \in \mathcal{P}_\tau(T')$, $P_1 \leq P_2$ if and only if $P_1 \subseteq P_2$, minimal elements of $\mathcal{P}_\tau(T')$ is denoted by

$$min\mathcal{P}_\tau(T') = \{P' \in \mathcal{P}_\tau(T') | \nexists P'' \in \mathcal{P}_\tau(T') \wedge P'' \subseteq P'\}.$$

Corollary 3. For $P' \in \min \mathcal{P}_\tau(T')$, P' is a minimal generator of $T'' \in [T]$, where T' is a generator of T'' .

In Example 1, $\{a_1, a_2, a_3\} \in \min[\{a_1, a_2, a_3\}]$, due to $\tau(\{a_1, a_2\}) = \tau(\{a_1, a_2, a_3\}) = \{a_5\}$, $\{a_1, a_2\} \in \min \mathcal{P}_\tau(\{a_1, a_2, a_3\})$ and $\{a_1, a_2\}$ is a minimal generator of $\{a_1, a_2, a_3\}$, $\{a_1, a_2, a_3, a_5\}$ and $\{a_1, a_2, a_3, a_4, a_5\}$.

From the algebraic point of view, we can construct a lattice structure on the quotient set of the topology $T(R_A)$, i.e., for any $[T_1], [T_2] \in T(R_A) / \sim_\tau$, we define $[T_1] \vee [T_2] = [(\cup[T_1]) \cap (\cup[T_2])]$ and $[T_1] \wedge [T_2] = [T_1 \cup T_2]$, one can easily check that operators \vee and \wedge on $T(R_A) / \sim_\tau$ are well defined, and we have the following theorem.

Theorem 3 ([52]). $(T(R_A) / \sim_\tau, \wedge, \vee)$ is a complete lattice, in which the maximum and minimum elements are $[\emptyset]$ and $[A]$, respectively.

In Example 1, $T(R_A) / \sim_\tau = \{[\emptyset], [\{a_1\}], [\{a_3\}], [\{a_5\}], [\{a_1, a_3\}], [\{a_1, a_5\}], [\{a_2, a_3\}], [\{a_3, a_5\}], [\{a_3, a_4, a_5\}], [\{a_1, a_2, a_3\}] (= [A])\}$, the lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$ of Example 1 is shown in Figure 2.

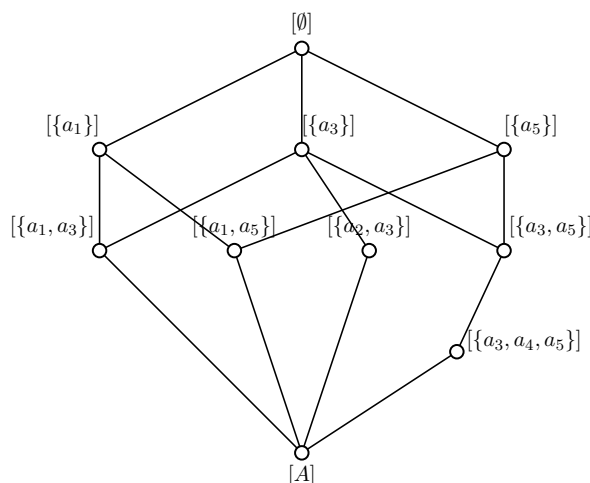


Figure 2. The complete lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$ of Example 1.

4. Association Rules Mining from the Quotient Set of the Topology

In summary, we can confirm the follows facts about itemsets of a transactional database $\mathcal{A} = (U, A)$ based on the above mentioned results:

1. The topology $T(R_A)$ for the set of items is a complete lattice and displays a hierarchical structure on some itemsets, it can be generated by the base $\mathcal{B}_A = \{C(a_i) | a_i \in A\}$; generally, each $C(a_i)$ is an itemset and can more fast generate closed itemsets than single items in the existed methods;
2. All closed itemsets are included in the topology $T(R_A)$, moreover, a closed itemset is the maximum element of an equivalent class $[T] \in T(R_A) / \sim_\tau$;
3. Each itemsets in $[T]$ has the same support; moreover, generators and minimal generators of a closed itemset can be obtained from $[T]$;
4. The complete lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$ displays the hierarchical structures on closed itemsets.

Inspired by rough set theory, each equivalent class in $T(R_A) / \sim_\tau$ can be understood as a granular knowledge on the set of itemsets of $\mathcal{A} = (U, A)$; on the one hand, each granular knowledge can be used to mine association rules because it includes a closed itemset. On the other hand, all granular knowledge can be also used to approximate any itemset $A_1 = \{a_{i_1}, a_{i_2}, \dots, a_{i_k}\} \subseteq A$, i.e.,

1. The lower approximation of A_1 :

$$L(A_1) = \bigcup_{a_i \in A}^{C(a_i) \subseteq A_1} C(a_i);$$

2. The upper approximation of A_1 :

$$G(A_1) = \bigcup_{a_i \in A}^{C(a_i) \cap A_1 \neq \emptyset} C(a_i).$$

It is obvious that for any itemset $A_1 \subseteq A$, $L(A_1)$ and $G(A_1)$ are in the topology $T(R_A)$ for the set of items, if $L(A_1) = G(A_1)$, then the itemset A_1 is in $T(R_A)$; furthermore, supports of $L(A_1)$, A_1 and $G(A_1)$ are such that

$$supp(L(A_1)) \geq supp(A_1) \geq supp(G(A_1)),$$

such as in Example 1, for itemset $\{a_1, a_2, a_5\} \subset A$, $L(\{a_1, a_2, a_5\}) = \{a_1, a_5\}$, $G(\{a_1, a_2, a_5\}) = \{a_1, a_2, a_3, a_5\}$, and $supp(\{a_1, a_5\}) = |\{u_1, u_5\}| = 2 \geq supp(\{a_1, a_2, a_5\}) = |\{u_5\}| = 1 \geq supp(\{a_1, a_2, a_3, a_5\}) = |\{u_5\}| = 1$. Formally, the lower approximation and the upper approximation of an itemset provide us an alternative method to mine generalized association rules or reduce association rules. All of these will be discussed in the rest of the section.

4.1. Min-Max Association Rules Mining

In the subsection, we provide an useful method to mine Min-Max association rules from closed frequent itemsets; here, an association rule is called as a Min-Max association rule if and only if there does not exist an association rule with the same quality measures as the association rule, but with a more specific antecedent part and a more general consequent part. Because closed itemsets are in equivalent classes of $T(R_A) / \sim_\tau$, we propose Algorithm 1 to generate all Min-Max association rules with confidence $c = 1$ from equivalent classes of $T(R_A) / \sim_\tau$.

Algorithm 1 Min-Max association rules mining from closed itemsets

Input: A transactional database $\mathcal{A} = (U, A)$.

Output: Min-Max association rules with confidence $c = 1$.

while The stop condition is not satisfied **do**

(1) Generate $\mathcal{B}_A = \{C(a_i) | a_i \in A\}$ according to $C(a_i) = \{a_j \in A | a_j \in \gamma\tau(\{a_i\})\}$ ($\forall a_i \in A$).

(2) Generate the topology $T(R_A) = \{\cup_{C(a_i) \in \mathcal{B}'_A} C(a_i) | \forall \mathcal{B}'_A \subseteq \mathcal{B}_A\}$ according to \mathcal{B}_A .

(3) Generate $T(R_A) / \sim_\tau = \{[T] | T \in T(R_A), \forall T_1, T_2 \in [T], \tau(T_1) = \tau(T_2)\}$ according to $T(R_A)$ and the mapping τ .

(4) Generate $min[T] = \{T' \in [T] | \nexists T'' \in [T] \wedge T'' \subseteq T'\}$ according to set inclusion for any $[T] \in T(R_A) / \sim_\tau$.

(5) Generate $P \in min\mathcal{P}_\tau(T')$ according to set inclusion for any $T' \in min[T]$.

(6) Select $P \in min\mathcal{P}_\tau(T')$ ($T' \in min[T]$), generate Min-Max association rule

$$\phi \equiv P \longrightarrow Q,$$

, in which $P \in min\mathcal{P}_\tau(T')$, $T' \in min[T]$, $Q = \cup[T] - P$ and $conf(\phi) = 1$.

(7) Return to (6).

end while

Output All $\phi \equiv P \longrightarrow Q$ with $conf(\phi) = 1$.

The pseudocode provided in Algorithm 1, is responsible for mining Min-Max association rules from closed itemsets. Step (1) generates the base \mathcal{B} , step (2) generates the topology $T(R_A)$, step (3) generates the equivalent class $T(R_A)/\sim_\tau$, step (4) generates the minimum of $[T]$.

Example 2. A transactional database is shown in Table 3. According to Definition 1, we obtain $C(a_1) = \{a_1, a_5\}$, $C(a_2) = \{a_2, a_5\}$, $C(a_3) = \{a_3\}$, $C(a_4) = \{a_3, a_4\}$ and $C(a_5) = \{a_5\}$, i.e., the base for the topology for the set of items is $\mathcal{B}_A = \{\{a_1, a_5\}, \{a_2, a_5\}, \{a_3\}, \{a_3, a_4\}, \{a_5\}\}$, the co-occurrence relation R_A among items decided by Equation (1) is shown in Table 4. The topology $T(R_A)$ for the set of items generated by \mathcal{B}_A and the quotient set $T(R_A)/\sim_\tau$ of the topology are

$$\begin{aligned}
 T(R_A) &= \{\emptyset, \{a_1, a_5\}, \{a_2, a_5\}, \{a_3\}, \{a_3, a_4\}, \{a_5\}, \{a_3, a_5\}, \{a_1, a_2, a_5\}, \\
 &\quad \{a_1, a_3, a_5\}, \{a_2, a_3, a_5\}, \{a_3, a_4, a_5\}, \{a_1, a_2, a_3, a_5\}, \{a_1, a_3, a_4, \\
 &\quad a_5\}, \{a_2, a_3, a_4, a_5\}, A(= \{a_1, a_2, a_3, a_4, a_5\})\}, \\
 T(R_A)/\sim_\tau &= \{\{\emptyset\}, [\{a_1, a_5\}], [\{a_2, a_5\}], [\{a_3\}], [\{a_3, a_4\}], [\{a_5\}], [\{a_3, a_5\}], \\
 &\quad [\{a_1, a_2, a_5\}], [\{a_2, a_3, a_5\}], [\{a_3, a_4, a_5\}], [A]\}.
 \end{aligned}$$

The complete lattice $(T(R_A)/\sim_\tau, \wedge, \vee)$ is shown in Figure 3, in which, each equivalent class is with its support, e.g., $([\{a_5\}], 7)$ means $\text{supp}(\{a_5\}) = 7$. Closed itemset and minimal generators of each equivalent class are shown in Table 5; accordingly, Min-Max association rules with confidence $c = 1$ generated from each equivalent class are shown in Table 6.

Table 3. A transactional database $\mathcal{A} = (U, A)$.

$U \setminus A$	a_1	a_2	a_3	a_4	a_5
u_1	1	0	1	0	1
u_2	0	0	1	1	0
u_3	0	1	0	0	1
u_4	1	0	1	0	1
u_5	0	0	1	1	0
u_6	1	1	1	0	1
u_7	1	0	1	1	1
u_8	1	1	0	0	1
u_9	1	0	0	0	1
u_{10}	0	0	1	1	0

Table 4. A binary relation on A .

R_A	a_1	a_2	a_3	a_4	a_5
a_1	1	0	0	0	1
a_2	0	1	0	0	1
a_3	0	0	1	0	0
a_4	0	0	1	1	0
a_5	0	0	0	0	1

Table 5. Closed itemset and minimal generators of each equivalent class.

Equivalent Class	$\cup[T]$	$\text{min}[T]$	$\text{min}\mathcal{P}_\tau(T')$
$\{a_1, a_5\}$	$\{a_1, a_5\}$	$\{a_1, a_5\}$	$\{a_1\}$
$\{a_2, a_5\}$	$\{a_2, a_5\}$	$\{a_2, a_5\}$	$\{a_2\}$
$\{a_3, a_4\}$	$\{a_3, a_4\}$	$\{a_3, a_4\}$	$\{a_4\}$
$\{a_1, a_2, a_5\}$	$\{a_1, a_2, a_5\}$	$\{a_1, a_2, a_5\}$	$\{a_1, a_2\}$
$\{a_3, a_5\}$	$\{a_1, a_3, a_5\}$	$\{a_3, a_5\}$	$\{a_3, a_5\}$
$\{a_2, a_3, a_5\}$	$\{a_1, a_2, a_3, a_5\}$	$\{a_2, a_3, a_5\}$	$\{a_2, a_3\}$
$\{a_3, a_4, a_5\}$	$\{a_1, a_3, a_4, a_5\}$	$\{a_3, a_4, a_5\}$	$\{a_4, a_5\}$

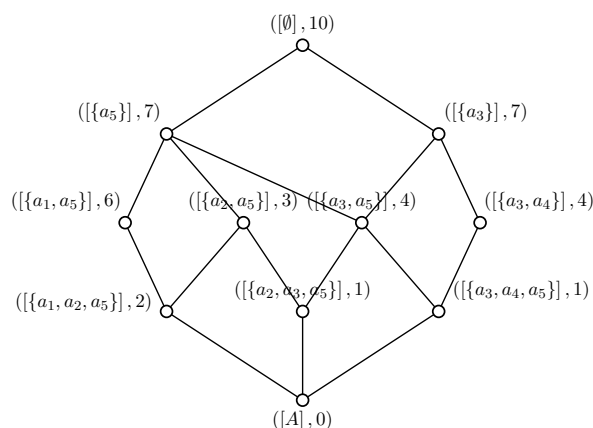


Figure 3. The complete lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$ of Example 2.

Table 6. Min-Max association rules with support s and confidence $c = 1$.

Min-Max Association Rule	(Support, Confidence)
$a_1 \rightarrow a_5$	(6, 1)
$a_2 \rightarrow a_5$	(3, 1)
$a_4 \rightarrow a_3$	(4, 1)
$a_1 \wedge a_2 \rightarrow a_5$	(2, 1)
$a_3 \wedge a_5 \rightarrow a_1$	(4, 1)
$a_2 \wedge a_3 \rightarrow a_5$	(1, 1)
$a_2 \wedge a_3 \wedge a_5 \rightarrow a_1$	(1, 1)
$a_2 \wedge a_3 \rightarrow a_1 \wedge a_5$	(1, 1)
$a_4 \wedge a_5 \rightarrow a_3$	(1, 1)
$a_3 \wedge a_4 \wedge a_5 \rightarrow a_1$	(1, 1)
$a_4 \wedge a_5 \rightarrow a_1 \wedge a_3$	(1, 1)

In Table 6, $a_2 \wedge a_3 \rightarrow a_5$ is generated by a minimal generator $\{a_2, a_3\}$ of $\{a_2, a_3, a_5\}$, $a_2 \wedge a_3 \wedge a_5 \rightarrow a_1$ is generated by generator $\{a_2, a_3, a_5\}$ of $\{a_1, a_2, a_3, a_5\}$, which are association rules with confidence $c = 1$; however, they are not Min-Max association rules with confidence $c = 1$, only $a_2 \wedge a_3 \rightarrow a_1 \wedge a_5$ is Min-Max association rules with confidence $c = 1$. Generally, there are many association rules with confidence $c = 1$ generated from each equivalent class; however, Min-Max association rules $P \rightarrow Q$ with confidence $c = 1$ must be generated by $P \in \min\mathcal{P}_\tau(T')$, $T' \in \min[T]$ and $Q = \cup[T] - P$, i.e., the advantage of our method is that searching minimal generator is limited in $T' \in \min[T]$ but not in all subsets of $\cup[T]$.

According to the hierarchical structures on closed itemsets displayed in the complete lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$, we provide Algorithm 2 to generate Min-Max association rules with high support and confidence from a fixed itemset.

The pseudocode provided in Algorithm 2 is responsible for mining Min-Max association rules from a fixed itemsets. Step (1) generates $L(A_1)$, step (2) generates the set of equivalent classes, and step (4) generates the Min-Max association rules.

In Example 2, several fixed itemsets and their lower approximations are shown in Table 7, the corresponding Min-Max association rules generated by the itemsets are shown in Table 8.

Algorithm 2 Min-Max association rules mining from a fixed itemset

Input: An itemset $A_1 = \{a_{i_1}, a_{i_2}, \dots, a_{i_k}\} \subseteq A$ with $supp(A_1) \geq s$.

Output: Min-Max association rules with confidence $c \in (0, 1)$.

while The stop condition is not satisfied **do**

(1) Generate $L(A_1)$ in $T(R_A)$.

(2) Generate the set of equivalent classes such that $A_1^{\sim\tau} = \{[T_i] \mid [T_i] \geq_{\sim\tau} [L(A_1)]\}$.

(3) Select $[T_i]$ from $A_1^{\sim\tau}$ such that $\cup[T_i] \subseteq A_1$,

(4) Generate association rule

$$\phi \equiv P \longrightarrow Q,$$

where, $P \in \min\mathcal{P}_\tau(T')$, $T' \in \min[T_i]$, $Q = A_1 - P$ and $conf(\phi) = \frac{supp(A_1)}{supp(T_i)}$.

(5) Return to (3).

end while

Output All $\phi \equiv P \longrightarrow Q$ with $supp(A_1) \geq s$ and $conf(\phi) = \frac{supp(A_1)}{supp(P)}$.

Table 7. Itemsets and their lower approximations of Example 2.

Itemsets	Lower Approximations	The Set of Equivalent Classes
$\{a_1, a_3\}$	$\{a_3\}$	$\{\{a_3\}\}$
$\{a_3, a_5\}$	$\{a_3, a_5\}$	$\{\{a_3\}, \{a_5\}\}$
$\{a_1, a_3, a_4\}$	$\{a_3, a_4\}$	$\{\{a_3\}, \{a_3, a_4\}\}$

Table 8. Min-Max association rules generated by the itemsets.

Itemsets	Min-Max Association Rules	(Support, Confidence)
$\{a_1, a_3\}$	$a_3 \longrightarrow a_1$	$(4, \frac{4}{7})$
$\{a_3, a_5\}$	$a_3 \longrightarrow a_5$	$(4, \frac{4}{7})$
$\{a_3, a_5\}$	$a_5 \longrightarrow a_3$	$(4, \frac{4}{7})$
$\{a_1, a_3, a_4\}$	$a_3 \longrightarrow a_1 \wedge a_4$	$(1, \frac{1}{7})$
$\{a_1, a_3, a_4\}$	$a_4 \longrightarrow a_1 \wedge a_3$	$(1, \frac{1}{4})$

Finally, advantages of Algorithms 1 and 2 can be summarized as follows:

1. Min-Max association rules are always mined from closed itemsets, in this paper, we prove that closed itemsets are maximum elements of equivalent classes, i.e., equivalent classes can be used to mine Min-Max association rules with confidence 1;
2. The shortest length antecedents of Min-Max association rules are searched from minimal members of equivalent classes, i.e., $P \in \min\mathcal{P}_\tau(T')$ and $T' \in \min[T]$; in this paper, searching minimal generators are in smaller scope than in all subsets of closed itemsets;
3. Lower approximations and their minimal generators help us to fast mine Min-Max association rules from a fixed itemset.

4.2. Generalized Association Rules Based on the Lower Approximation

Generalized association rules are important extension of association rules. By using taxonomy over items of transactional databases, generalized items are aggregated according to different granularity levels. In practice, generalized itemsets provide a high level view of the patterns hidden in the analyzed data and a high level abstraction of the mined knowledge in different application domains [12,30,59,60]. In this paper, each equivalent class $[T]$ in $T(R_A) / \sim_\tau$ is understood as a granular knowledge on the set of itemsets; it can also be considered as a kind of generalized itemsets. Accordingly, we provide an alternative method to mine generalized association rules based on the lower approximation of any itemset, and analyze the changing confidence of generalized association rules. Formally, for any association rule $\psi \equiv A_1 \longrightarrow A_2$, we propose the following three kinds of generalized association rules of ψ based on the lower approximation:

1. Generalized antecedent association rule (GAR): $\psi_{GA} \equiv L(A_1) \longrightarrow A_2$;
2. Generalized conclusion association rule (GCR): $\psi_{GC} \equiv A_1 \longrightarrow L(A_2)$;
3. Generalized antecedent and conclusion association rule (GACR): $\psi_{GAC} \equiv L(A_1) \longrightarrow L(A_2)$.

It is obvious that if $A_1 \in T(R_A)$, then $L(A_1) = A_1$ and $\psi = \psi_{GA}$; if $A_2 \in T(R_A)$, then $L(A_2) = A_2$ and $\psi = \psi_{GC}$; if $A_1, A_2 \in T(R_A)$, then $L(A_1) = A_1, L(A_2) = A_2$ and $\psi = \psi_{GAC}$.

Corollary 4. For any association rule $\psi \equiv A_1 \longrightarrow A_2$, $conf_{\mathcal{S}}(\psi) = c, l = |\tau(L(A_1)) - \tau(A_1)|$ and $k = |(\tau(L(A_1)) - \tau(A_1)) \cap \tau(A_2)|$.

1. $conf_{\mathcal{S}}(\psi) \leq conf_{\mathcal{S}}(\psi_{GC})$;
2. $conf_{\mathcal{S}}(\psi_{GA}) \leq conf_{\mathcal{S}}(\psi_{GAC})$;
3. If $k \geq cl$, then $conf_{\mathcal{S}}(\psi) \leq conf_{\mathcal{S}}(\psi_{GA})$;
4. If $k < cl$, then $conf_{\mathcal{S}}(\psi) > conf_{\mathcal{S}}(\psi_{GA})$.

Proof. Based on $\tau(A_1) \subseteq \tau(L(A_1))$ and $\tau(A_2) \subseteq \tau(L(A_2))$, (1) and (2) can be easily proved. For (3) and (4), due to $\tau(L(A_1)) = \tau(A_1) \cup (\tau(L(A_1)) - \tau(A_1))$ and $\tau(A_1) \cap (\tau(L(A_1)) - \tau(A_1)) = \emptyset$, we have $conf_{\mathcal{S}}(\psi_{GA}) = \frac{|\tau(L(A_1)) \cap \tau(A_2)|}{|\tau(L(A_1))|} = \frac{|(\tau(A_1) \cup (\tau(L(A_1)) - \tau(A_1))) \cap \tau(A_2)|}{|\tau(L(A_1))|} = \frac{|(\tau(A_1) \cap \tau(A_2)) \cup ((\tau(L(A_1)) - \tau(A_1)) \cap \tau(A_2))|}{|\tau(L(A_1))|} = \frac{k + |(\tau(L(A_1)) - \tau(A_1)) \cap \tau(A_2)|}{l + |\tau(A_1)|}$, by $\frac{|\tau(A_1) \cap \tau(A_2)|}{|\tau(A_1)|} = c$, if $k \geq cl$, then $\frac{k + |(\tau(L(A_1)) - \tau(A_1)) \cap \tau(A_2)|}{l + |\tau(A_1)|} \geq \frac{cl + c|\tau(A_1)|}{l + |\tau(A_1)|} = c$. If $k < cl$, then $\frac{k + |(\tau(L(A_1)) - \tau(A_1)) \cap \tau(A_2)|}{l + |\tau(A_1)|} < \frac{cl + c|\tau(A_1)|}{l + |\tau(A_1)|} = c$. \square

Corollary 5. Let $conf_{\mathcal{S}}(\psi_{GC}) = c, l = |\tau(L(A_1)) - \tau(A_1)|$ and $k = |(\tau(L(A_1)) - \tau(A_1)) \cap \tau(L(A_2))|$. 1) If $k \geq cl$, then $conf_{\mathcal{S}}(\psi_{GC}) \leq conf_{\mathcal{S}}(\psi_{GAC})$; 2) If $k < cl$, then $conf_{\mathcal{S}}(\psi_{GC}) > conf_{\mathcal{S}}(\psi_{GAC})$.

Based on $P \in \min \mathcal{P}_{\tau}(T'), T' \in \min[L(A_1)]$ and $\cup[L(A_2)]$, ψ_{GA}, ψ_{GC} and ψ_{GAC} of association rule $\psi \equiv A_1 \longrightarrow A_2$ have the following redundant rules with the same confidence and the same support:

1. Redundant association rule of GAR: $\psi_{GAN} \equiv P \longrightarrow A_2$;
2. Redundant association rule of GCR: $\psi_{GCN} \equiv A_1 \longrightarrow \cup[L(A_2)]$;
3. Redundant association rule of GACR: $\psi_{GACN} \equiv P \longrightarrow \cup[L(A_2)]$.

Accordingly, we provide Algorithm 3 to mine generalized association rules of any association rule ψ and its corresponding redundant association rules.

Algorithm 3 Mining generalized association rules and redundant association rules based on the lower approximation

Input: Association rule $\psi \equiv A_1 \longrightarrow A_2$.

Output: Generalized association rules of ψ and corresponding redundant association rules.

while The stop condition is not satisfied **do**

(1) Generate $L(A_1)$ and $L(A_2)$ in $T(R_A)$.

(2) Generate $\cup[L(A_1)]$ and $\cup[L(A_2)]$ according to set inclusion of $[L(A_1)]$ and $[L(A_2)]$, respectively.

(3) Generate $T' \in \min[L(A_1)]$ and $P \in \min \mathcal{P}_{\tau}(T')$ according to set inclusion of $[L(A_1)]$ and the power set of T' , respectively.

(4) Obtain generalized association rules $\psi_{GA} \equiv L(A_1) \longrightarrow A_2, \psi_{GC} \equiv A_1 \longrightarrow L(A_2)$ and $\psi_{GAC} \equiv L(A_1) \longrightarrow L(A_2)$, respectively.

(5) Generate redundant association rules $\psi_{GAN} \equiv P \longrightarrow A_2, \psi_{GCN} \equiv A_1 \longrightarrow \cup[L(A_2)]$ and $\psi_{GACN} \equiv P \longrightarrow \cup[L(A_2)]$, respectively.

(6) Return to (5).

end while

Output Generalized association rules and corresponding redundant association rules based on the lower approximation with $(supp(\psi_*), conf(\psi_*))$.

The pseudocode provided in Algorithm 3 is responsible for mining generalized association rules and redundant association rules based on the lower approximation. Step (1) generates $L(A_1)$ and $L(A_2)$, step (2) generates $\cup[L(A_1)]$ and $\cup[L(A_2)]$, step (4) generates association rules.

4.3. Generalized Association Rules Based on the Upper Approximation

Similarly, for any association rule $\psi \equiv A_1 \rightarrow A_2$, we propose the following three kinds of generalized association rules of ψ based on the upper approximation:

1. Generalized antecedent association rule (gar): $\psi_{ga} \equiv G(A_1) \rightarrow A_2$;
2. Generalized conclusion association rule (gcr): $\psi_{gc} \equiv A_1 \rightarrow G(A_2)$;
3. Generalized antecedent and conclusion association rule (gacr): $\psi_{gac} \equiv G(A_1) \rightarrow G(A_2)$.

It is obvious that if $A_1 \in T(R_A)$, then $\psi = \psi_{ga}$; if $A_2 \in T(R_A)$, then $\psi = \psi_{gc}$; if $A_1, A_2 \in T(R_A)$, and then $\psi = \psi_{gac}$.

Corollary 6. For $\psi \equiv A_1 \rightarrow A_2$, $conf_{\mathcal{A}}(\psi) = c$, $l = |\tau(A_1) - \tau(G(A_1))|$ and $k = |(\tau(A_1) - \tau(G(A_1))) \cap \tau(A_2)|$.

1. $conf_{\mathcal{A}}(\psi) \geq conf_{\mathcal{A}}(\psi_{gc})$;
2. $conf_{\mathcal{A}}(\psi_{ga}) \geq conf_{\mathcal{A}}(\psi_{gac})$;
3. If $k > cl$, then $conf_{\mathcal{A}}(\psi) > conf_{\mathcal{A}}(\psi_{ga})$;
4. If $k \leq cl$, then $conf_{\mathcal{A}}(\psi) \leq conf_{\mathcal{A}}(\psi_{ga})$.

Proof. Based on $\tau(A_1) \supseteq \tau(G(A_1))$ and $\tau(A_2) \supseteq \tau(G(A_2))$, (1) and (2) can be easily proved. For (3) and (4), due to $\tau(A_1) = \tau(G(A_1)) \cup (\tau(A_1) - \tau(G(A_1)))$ and $\tau(G(A_1)) \cap (\tau(A_1) - \tau(G(A_1))) = \emptyset$, we have

$$\begin{aligned} conf_{\mathcal{A}}(\psi) &= \frac{|\tau(A_1) \cap \tau(A_2)|}{|\tau(A_1)|} = \frac{|\tau(G(A_1)) \cup (\tau(A_1) - \tau(G(A_1))) \cap \tau(A_2)|}{|\tau(G(A_1)) \cup (\tau(A_1) - \tau(G(A_1)))|} \\ &= \frac{|\tau(G(A_1)) \cap \tau(A_2) \cup ((\tau(A_1) - \tau(G(A_1))) \cap \tau(A_2))|}{|\tau(G(A_1)) \cup (\tau(A_1) - \tau(G(A_1)))|} \\ &= \frac{|\tau(G(A_1)) \cap \tau(A_2)| + k}{|\tau(G(A_1))| + l} = c, \end{aligned}$$

i.e., $|\tau(G(A_1)) \cap \tau(A_2)| - c|\tau(G(A_1))| = cl - k$. If $k > cl$, then $|\tau(G(A_1)) \cap \tau(A_2)| - c|\tau(G(A_1))| < 0$, i.e., $conf_{\mathcal{A}}(\psi_{ga}) = \frac{|\tau(G(A_1)) \cap \tau(A_2)|}{|\tau(G(A_1))|} < c$. If $k \leq cl$, then $|\tau(G(A_1)) \cap \tau(A_2)| - c|\tau(G(A_1))| \geq 0$, i.e., $conf_{\mathcal{A}}(\psi_{ga}) = \frac{|\tau(G(A_1)) \cap \tau(A_2)|}{|\tau(G(A_1))|} \geq c$. \square

Corollary 7. Let $conf_{\mathcal{A}}(\psi_{gc}) = c$, $l = |\tau(A_1) - \tau(G(A_1))|$ and $k = |(\tau(A_1) - \tau(G(A_1))) \cap \tau(G(A_2))|$. (1) If $k > cl$, then $conf_{\mathcal{A}}(\psi_{gc}) > conf_{\mathcal{A}}(\psi_{gac})$; (2) If $k \leq cl$, then $conf_{\mathcal{A}}(\psi_{gc}) \leq conf_{\mathcal{A}}(\psi_{gac})$.

Based on $P \in \min \mathcal{P}_{\tau}(T')$, $T' \in \min[G(A_1)]$ and $\cup[G(A_2)]$, ψ_{ga} , ψ_{gc} and ψ_{gac} of association rule $\psi \equiv A_1 \rightarrow A_2$ have the following redundant rules with the same confidence and support:

1. Redundant association rule of gar: $\psi_{gan} \equiv P \rightarrow A_2$;
2. Redundant association rule of gcr: $\psi_{gcn} \equiv A_1 \rightarrow \cup[G(A_2)]$;
3. Redundant association rule of gacr: $\psi_{gacn} \equiv P \rightarrow \cup[G(A_2)]$.

Accordingly, we provide Algorithm 4 to mine generalized association rules of ψ and corresponding redundant association rules.

The pseudocode provided in Algorithm 4, is responsible for mining generalized association rules and redundant association rules based on the upper approximation. Step (1) generates $G(A_1)$ and $G(A_2)$, step (2) generates $\cup[G(A_1)]$ and $\cup[G(A_2)]$, step (4) generates association rules, and step (5) generates redundant association rules.

Algorithm 4 Mining generalized association rules and redundant association rules based on the upper approximation

Input: Association rule $\psi \equiv A_1 \rightarrow A_2$.

Output: Generalized association rules of ψ and corresponding redundant association rules.

while The stop condition is not satisfied **do**

(1) Generate $G(A_1)$ and $G(A_2)$ in $T(R_A)$.

(2) Generate $\cup[G(A_1)]$ and $\cup[G(A_2)]$ according to set inclusion of $[G(A_1)]$ and $[G(A_2)]$, respectively.

(3) Generate $T' \in \min[G(A_1)]$ and $P \in \min\mathcal{P}_\tau(T')$ according to set inclusion of $[G(A_1)]$ and the power set of T' , respectively.

(4) Obtain generalized association rules $\psi_{gar} \equiv G(A_1) \rightarrow A_2$, $\psi_{gcr} \equiv A_1 \rightarrow G(A_2)$ and $\psi_{gacr} \equiv G(A_1) \rightarrow G(A_2)$, respectively.

(5) Generate redundant association rules $\psi_{gan} \equiv P \rightarrow A_2$, $\psi_{gcn} \equiv A_1 \rightarrow \cup[G(A_2)]$ and $\psi_{gacn} \equiv P \rightarrow \cup[G(A_2)]$, respectively.

(6) Return to (5).

end while

Output Generalized association rules and corresponding redundant association rules based on the upper approximation with $(supp(\psi_*), conf(\psi_*))$.

5. Example Analysis

Experiments were made to compare the execution time, memory usage and numbers of association rules of the Apriori [1,61] algorithms and our method. They were implemented on a Thinkpad X1 laptop with Intel i5 Core Duo (2 × 2.4 GHz), 4 GB of RAM and running Windows 10. The algorithms were coded in Matlab 2015b. Four databases from the UCI databases [62] were used for the experiments, of which the features are shown in Table 9.

5.1. The Execution Time

Experiments were made to compare the execution time of the algorithms Apriori and ours. The minConf was set to 50%. The results of the four databases for various minSup values are shown in Figures 4–7, which shows that the execution of our algorithm was faster than Apriori in all cases. For example, given $minSup = 75\%$, for the Chess database, the mining time of Apriori and ours were 3035.53(s) and 507.33(s), respectively. The time ratio is $\frac{507.33}{3035.53} \times 100\% = 16.71\%$. Besides, as the minSup is decreased, the time ratio is reduced also. For example, consider the Connect database with $minSup$ set at 98%, 96% and 94%, the speed up of the time ratio were 91.93%, 7.26% and 1.56%, respectively. The results demonstrate that the execution time of mining rules from ours was more efficient than Apriori.

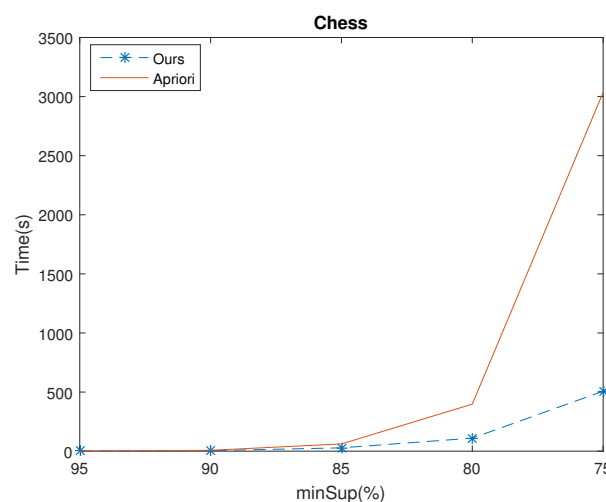


Figure 4. Execution time of the two algorithms in Chess dataset for various minSup values.

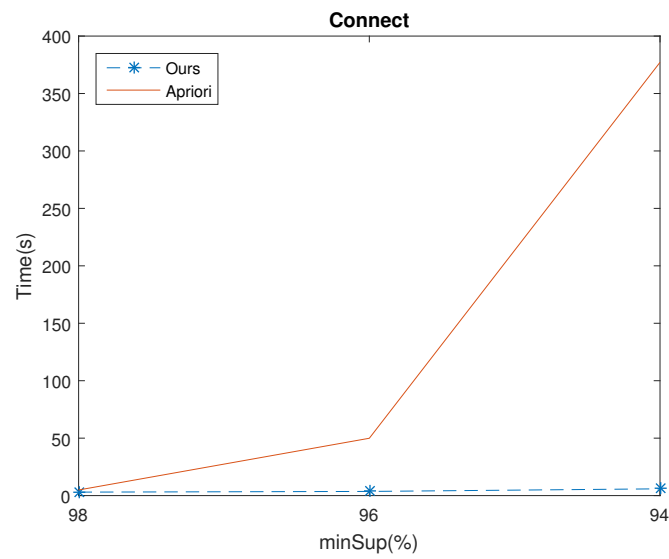


Figure 5. Execution time of the two algorithms in Connect dataset for various minSup values.

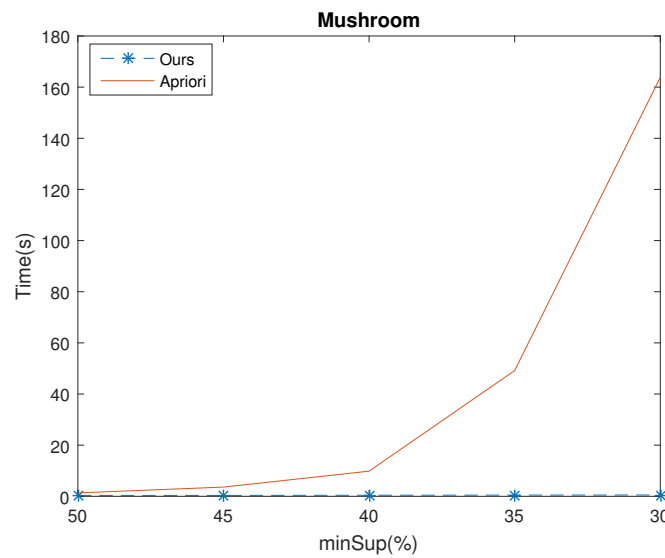


Figure 6. Execution time of the two algorithms in Mushroom dataset for various minSup values.

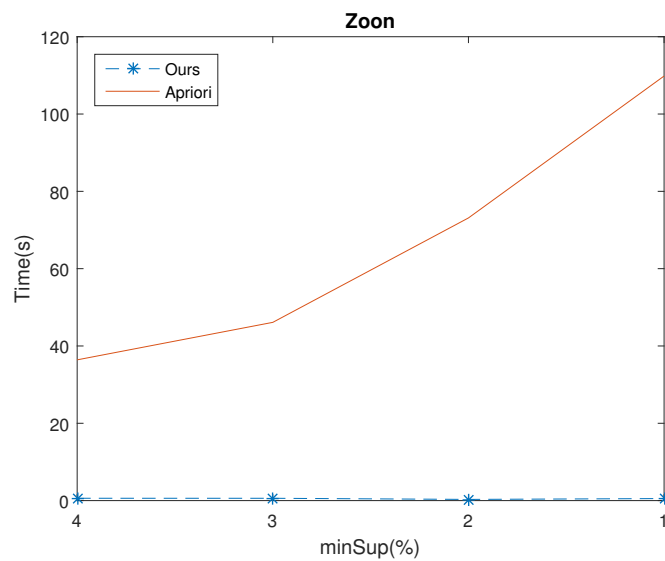


Figure 7. Execution time of the two algorithms in Zoon dataset for various minSup values.

5.2. The Memory Usage

Experiments were made with the same databases and same parameters as the execution time experiments. The results show that the Apriori algorithm consumed more memory than ours in almost all cases, which are shown in Figures 8–11. For example, given $minSup = 75\%$, for Chess database, the memory usage of Apriori and ours were 672 (Mb) and 136.54 (Mb), respectively. The time ratio is $\frac{136.54}{672} \times 100\% = 20.32\%$. Moreover, when we decrease the $minSup$, the time ratio will reduce as well. For example, consider the Connect database with $minSup$ set at 98%, 96% and 94%, the speed up of the time ratio were 87.55%, 75.84% and 9.18%, respectively.

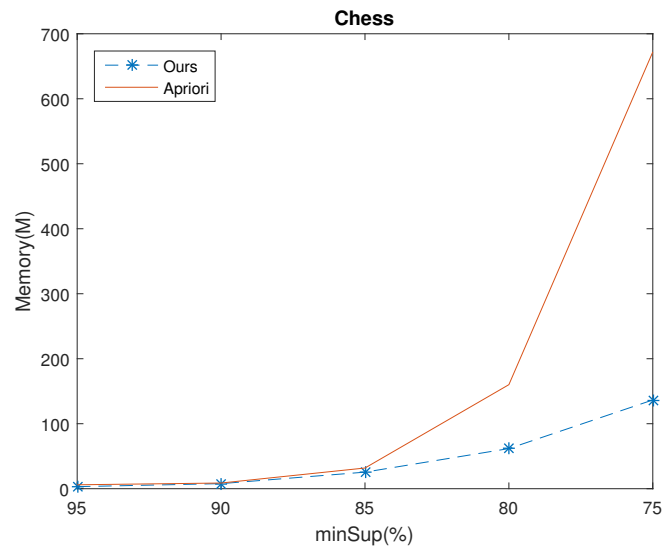


Figure 8. Memory usage of the two algorithms in Chess dataset for various minSup values.

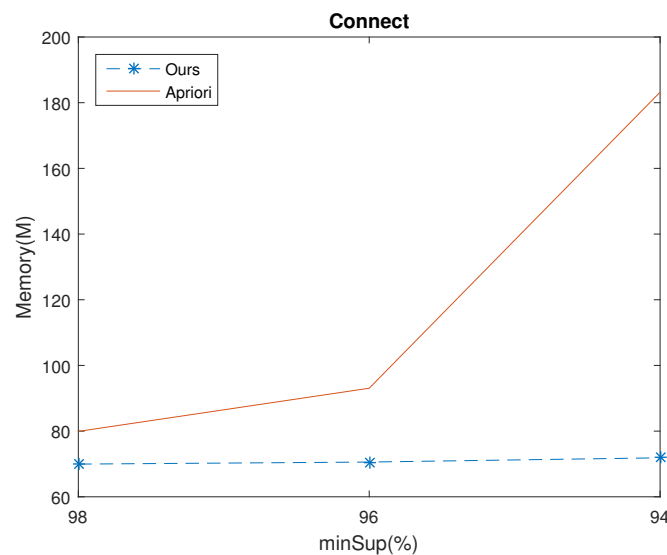


Figure 9. Memory usage of the two algorithms in Connect dataset for various minSup values.

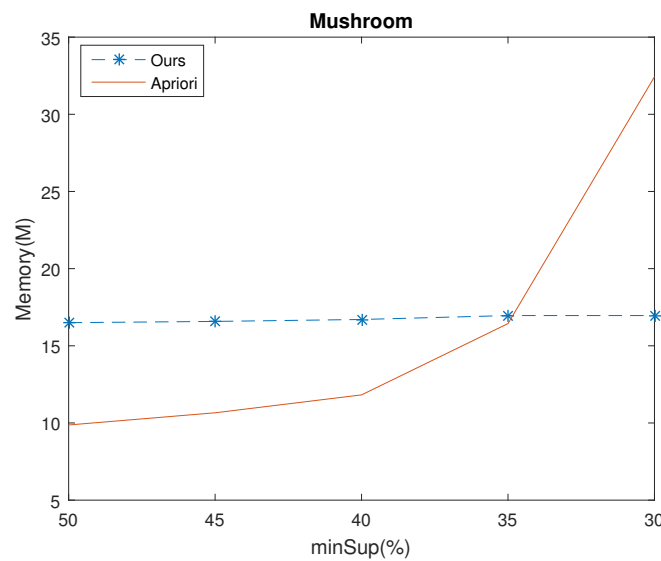


Figure 10. Memory usage of the two algorithms in Mushroom dataset for various minSup values.

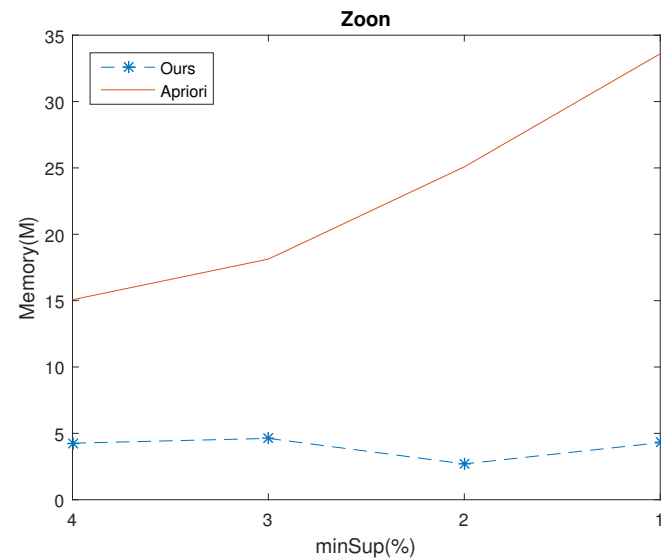


Figure 11. Memory usage of the two algorithms in Zoon dataset for various minSup values.

5.3. Numbers of Rules

We compare the numbers of association rules of Apriori with ours with the same databases and same parameters as the execution time experiments in Table 10. The results demonstrate that the numbers of ours is always smaller than those of Apriori. For example, given $minSup = 75%$, for Chess database, the number of rules of Apriori and ours were 2,336,556 and 253,836, respectively. The time ratio is $\frac{253836}{2336556} \times 100\% = 10.86\%$. For our algorithm generating Min-Max association rules, the numbers of rules of ours is smaller than Apriori.

Table 9. Dataset characteristics.

Dataset	Transactions	Original Attributes	Attributes after Conversion
Zoo	101	17	15
Mushroom	8124	23	126
Connect-4	67,557	43	129
Chess	3196	36	108

Table 10. Number of rules of Apriori and Ours.

Dataset	minSup(%)	Number of Rules	
		Apriori	Ours
Chess	95	472	700
	90	10,742	9482
	85	95,482	43,116
	80	552,564	111,768
	75	2,336,556	253,836
Mushroom	50	1146	172
	45	2704	291
	40	5006	483
	35	14,107	903
	30	45,145	903
Connect	98	1544	380
	96	27,340	1480
	94	201,928	3848
Zoon	4	29,288	8136
	3	35,204	8826
	2	48,578	5110
	1	64,868	8174

In this section, some examples are provided to show our method in association rules mining; transactional databases come from [62] and their characteristics are shown in Table 9. The Zoo database is initially used for classification of animals; in this paper, association rules among attributes of animals are considered, in which Min-Max association rules with confidence $c = 1$ and confidence $c \in (0, 1)$ will be shown, respectively. For mushroom database, connect-4 and chess, the topology, basis for the topology and equivalent classes with thresholds of support are obtained, and Min-Max association rules and generalized association rules are generated. Because the Reliable basis used to retrieve and reduce association rules in [63] is similar to our method, our results are compared with the method based on the Reliable basis limited in the number of Min-Max association rules and redundant rules.

Example 3. In the Zoo database, there are 101 objects (animals) and 17 attributes (15 boolean, 2 numeric). Here, we discover association rules among 15 boolean attributes of animals, i.e., hair (ha), feathers (fe), eggs (eg), milk (mi), airborne (ai), aquatic (aq), predator (pr), toothed (to), backbone (ba), breathes (br), venomous (ve), fins (fi), tail (ta), domestic (do) and catsize (ca).

According to Definition 1, 15 bases are generated (shown in Table A1 of Appendix A). Members of topology and equivalent classes with $supp_r(*)$ are shown in Table A2 of Appendix A. Equivalent classes used to generate association rules with confidence $c = 1$ are shown in Table A3 of Appendix A. According to Table A3, 7 Min-Max association rules with $supp_r(*) \geq 0.4$ and confidence $c = 1$ are generated (shown in Table A4 of Appendix A), in which $mi \rightarrow ba \wedge br$ (rule 4) is considered redundant of $mi \wedge ba \rightarrow br$ (rule 5) and $mi \wedge br \rightarrow ba$ (rule 6) due to $\{mi\} \subset \{mi, ba\}$, $\{mi\} \subset \{mi, br\}$, $\{br\} \subset \{ba, br\}$ and $\{ba\} \subset \{ba, br\}$; i.e., we use less conditions to achieve more results by $mi \rightarrow ba \wedge br$ with the same support and confidence. Furthermore, some Min-Max association rules generated in the quotient lattice of the topology with $supp_r(*) \geq 0.4$ and confidence $c \in (0, 1)$ are shown in Table A5 of Appendix A, in which the Min-Max association rule $br \rightarrow ba$ and $ta \rightarrow ba$ can be considered as a GAR of $br \wedge ta \rightarrow ba$, respectively.

Example 4. The mushroom database consists of a database with 8124 objects (mushrooms) and 22 nominally valued attributes. Here, we convert 22 nominally valued attributes as 126 boolean attributes and generate 111 bases. Change of bases with different $supp_r(*)$ is shown in Figure A1 of Appendix A. Members of topology, equivalent classes and its generating time by using different bases are shown in Figure A2 of Appendix A.

Table A6 of Appendix A shows Min-Max association rules with $\text{supp}_r(*) \geq 0.8$ and $c = 1$. Compared with results in [63], rules 1–8 and 11 are same with [63], and rules 9, 10, 12 and 13 are new rules, in which rule 8 can be considered as a redundant rule of rules 9 and 10, rule 11 being a redundant rule of rules 12 and 13.

Table A7 of Appendix A show Min-Max association rules with $\text{supp}_r(*) \geq 0.8$ and $c \in [0.8, 1)$, in which rules 1–25 are also generated in [63]; rules 26–41 are new rules.

Table A8 of Appendix A shows comparative results, in which TRX means the number of all approximate rules in [63], Min-Max (N, Rt) and Reliable (N, Rt) means number and reduction ratio of Min-Max basis rules and reliable basis rules generated by Min-Max approximate basis, Min-Max exact basis, reliable approximate basis and reliable exact basis in [63]. TR means the number of all rules generated in the quotient lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$, N-Min-Max (N, Rt) means the number and reduction ratio of Min-Max association rules generated in $(T(R_A) / \sim_\tau, \wedge, \vee)$, the reduction ratio is calculated as $Rt = \frac{TRX-N}{TRX}$ or $\frac{TR-N}{TR}$. In Table A8, it can be noticed that the association rules used to obtain a reduction in the quotient lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$ are less than in [63], such as for $\text{supp}_r(*) \geq 0.4$, 2528 association rules are used to obtain reduction. The lower $\text{supp}_r(*)$ and confidence c is corresponding to more generated association rules and higher reduction ratio in the quotient lattice $(T(R_A) / \sim_\tau, \wedge, \vee)$, such as for $\text{supp}_r(*) \geq 0.4$ and confidence $c \geq 0.5$, 1825 rules are generated and the reduction ratio is 0.77.

Example 5. In the connect-4 database, the number of reduction association rules with confidence $c \geq 0.5$ and compared with in [63] are shown in Table A9 of Appendix A. For template $A = \{d_5 - b, d_6 - b, e_5 - b, f_5 - b, f_6 - b\}$ and rule $\phi \equiv a_6 - b \wedge e_4 - b \wedge f_5 - b \rightarrow e_6 - b \wedge f_6 - b$,

$$\begin{aligned} L(\{d_5 - b, d_6 - b, e_5 - b, f_5 - b, f_6 - b\}) &= \{d_5 - b, d_6 - b, f_5 - b, f_6 - b\}, \\ \min[\{d_5 - b, d_6 - b, f_5 - b, f_6 - b\}] &= \{d_5 - b, f_5 - b\}, \\ L(\{a_6 - b, e_4 - b, f_5 - b\}) &= \{a_6 - b\}, \\ G(\{a_6 - b, e_4 - b, f_5 - b\}) &= \{a_6 - b, e_4 - b, e_5 - b, f_5 - b\}, \\ L(\{e_6 - b, f_6 - b\}) &= \{e_6 - b, f_6 - b\}, \\ G(\{e_6 - b, f_6 - b\}) &= \{e_6 - b, f_6 - b\}. \end{aligned}$$

Accordingly, Table A10 shows the association rule of the template and generalized association rules in the connect-4 database.

In the chess database, the number of reduction association rules with confidence $c \geq 0.5$ and compared with in [63] are shown in Table A11 of Appendix A. For template $A = \{\text{mulch}_-f, \text{skach}_-f, \text{wkna8}_-f\}$ and rule $\phi \equiv \text{qxmsq}_-f \wedge \text{spcop}_-f \wedge \text{wkna8}_-f \rightarrow \text{bkon8}_-f \wedge \text{hdchk}_-f$,

$$\begin{aligned} L(\{\text{mulch}_-f, \text{skach}_-f, \text{wkna8}_-f\}) &= \{\text{mulch}_-f, \text{skach}_-f\}, \\ \min[\{\text{mulch}_-f, \text{skach}_-f\}] &= \{\text{mulch}_-f, \text{skach}_-f\}, \\ L(\{\text{qxmsq}_-f, \text{spcop}_-f, \text{wkna8}_-f\}) &= \{\text{qxmsq}_-f, \text{spcop}_-f\}, \\ G(\{\text{qxmsq}_-f, \text{spcop}_-f, \text{wkna8}_-f\}) &= \{\text{qxmsq}_-f, \text{spcop}_-f, \text{wkna8}_-f, \text{stlmt}_-f\}, \\ L(\{\text{bkon8}_-f, \text{hdchk}_-f\}) &= \{\text{bkon8}_-f, \text{hdchk}_-f, \text{reskd}_-f\}, \\ G(\{\text{bkon8}_-f, \text{hdchk}_-f\}) &= \{\text{bkon8}_-f, \text{hdchk}_-f, \text{reskd}_-f, \text{thrsk}_-f, \text{spcop}_-f\}. \end{aligned}$$

Accordingly, Table A12 shows the association rule of the template and generalized association rules in the chess database.

6. Conclusions

Association rules are often generated by frequent itemsets or closed itemsets from transactional databases. In order to obtain these itemsets, many methods have been proposed. In this article, for representing more general associative relations, which are among the items of transaction databases, the topology on itemset of a transactional database can be constructed. The topology on itemset includes frequent itemsets and closed itemsets. The topology on itemset includes frequent itemsets and closed itemsets, which has been proved.

Most important of all, the basis of the topology can be used to generate the topology on the itemset, which deduced from the transactional database. Using the basis of the topology can efficiently avoid scanning databases many times in extracting association rules. The quotient lattice of the topology displays the hierarchical structures on all itemsets, because every closed itemset and its generators or minimal generators are limited in an element of the quotient lattice, valid Min-Max association rules can be easily generated in the element. Because the quotient lattice of the topology provides granular concepts to approximate any template of itemset, reductant association rules can be easily generated by granular concepts. The experiment demonstrates that association rules mining using topology for itemset is an efficient method.

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Data Availability Statement: Not applicable.

Conflicts of Interest: The authors declare no conflict of interest.

Appendix A

Table A1. Basis \mathcal{B}_A of Zoo database.

Attributes	<i>ha</i>	<i>fe</i>	<i>eg</i>	<i>mi</i>	<i>ai</i>
Basis	{ <i>ha, br</i> }	{ <i>fe, eg, ba, br, ta</i> }	{ <i>eg</i> }	{ <i>mi, ba, br</i> }	{ <i>ai, br</i> }
Attributes	<i>aq</i>	<i>pr</i>	<i>to</i>	<i>ba</i>	<i>br</i>
Basis	{ <i>aq</i> }	{ <i>pr</i> }	{ <i>to, ba</i> }	{ <i>ba</i> }	{ <i>br</i> }
Attributes	<i>ve</i>	<i>fi</i>	<i>ta</i>	<i>do</i>	<i>ca</i>
Basis	{ <i>ve</i> }	{ <i>aq, to, ba, fi</i> }	{ <i>ta</i> }	{ <i>do</i> }	{ <i>ca</i> }

Table A2. Members of topology and equivalent classes with $supp_r(*) \geq c$.

<i>c</i>	0	0.2	0.3	0.5
Members of topology	4159	98	54	11
Equivalent classes	237	75	44	11

Table A3. Equivalent classes used to generate association rules with $supp_r(*) \geq 0.3$ and confidence $c = 1$.

Equivalent Classes	$ [*] $	$min[*]$	$\mathcal{P}_\tau(T')$
$\{\{ba, br, ca\}\}$	2	$\{\{br, ca\}\}$	$\{\{br, ca\}\}$
$\{\{ba, br, ta, ca\}\}$	2	$\{\{br, ta, ca\}\}$	$\{\{br, ta, ca\}\}$
$\{\{ba, br, to\}\}$	1	$\{\{ba, br, to\}\}$	$\{\{br, to\}\}$
$\{\{ha, br\}\}$	1	$\{\{ha, br\}\}$	$\{\{ha\}\}$
$\{\{to, ba, br, ta\}\}$	1	$\{\{to, ba, br, ta\}\}$	$\{\{br, to, ta\}\}$
$\{\{mi, to, ba, br\}\}$	1	$\{\{mi, to, ba, br\}\}$	$\{\{mi, to\}, \{mi, to, ba\}, \{mi, to, br\}\}$
$\{\{ha, mi, ba, br\}\}$	2	$\{\{ha, ba, br\}\}$	$\{\{ha, ba\}, \{ha, ba, br\}\}$
$\{\{mi, to, ba, br, ta\}\}$	1	$\{\{mi, to, ba, br, ta\}\}$	$\{\{mi, to, ta\}, \{mi, to, ba, ta\}, \{mi, to, br, ta\}\}$
$\{\{eg, ba, ta\}\}$	2	$\{\{eg, ta\}\}$	$\{\{eg, ta\}\}$
$\{\{ha, mi, ba, br, ta\}\}$	3	$\{\{ha, br, ta\}\}$	$\{\{ha, ta\}, \{ha, br, ta\}\}$
$\{\{ha, mi, to, ba, br, ta\}\}$	2	$\{\{ha, to, ba, br, ta\}\}$	$\{\{ha, to, ta\}, \{ha, to, ba, br, ta\}, \{ha, to, br, ta\}, \{ha, to, ba, ta\}\}$
$\{\{pr, to, ba, ta\}\}$	1	$\{\{pr, to, ba, ta\}\}$	$\{\{pr, to, ta\}\}$
$\{\{to, ba\}\}$	1	$\{\{to, ba\}\}$	$\{\{to\}\}$
$\{\{pr, to, ba\}\}$	1	$\{\{pr, to, ba\}\}$	$\{\{pr, to\}\}$
$\{\{mi, ba, br\}\}$	1	$\{\{mi, ba, br\}\}$	$\{\{mi\}, \{mi, ba\}, \{mi, br\}\}$
$\{\{mi, ba, br, ta\}\}$	1	$\{\{mi, ba, br, ta\}\}$	$\{\{mi, ta\}, \{mi, ba, ta\}, \{mi, br, ta\}\}$
$\{\{to, ba, ca\}\}$	1	$\{\{to, ba, ca\}\}$	$\{\{to, ca\}\}$
$\{\{mi, ba, br, ca\}\}$	1	$\{\{mi, ba, br, ca\}\}$	$\{\{mi, ca\}, \{mi, ba, ca\}, \{mi, br, ca\}\}$
$\{\{mi, to, ba, br, ca\}\}$	2	$\{\{to, ba, br, ca\}\}$	$\{\{to, ba, ca\}, \{to, ba, br, ca\}\}$
$\{\{to, ba, ta\}\}$	1	$\{\{to, ba, ta\}\}$	$\{\{to, ta\}\}$
$\{\{ha, mi, to, ba, br\}\}$	2	$\{\{ha, to, ba, br\}\}$	$\{\{ha, to\}, \{ha, to, ba\}, \{ha, to, br\}, \{ha, to, ba, br\}\}$
$\{\{ba, ta, ca\}\}$	2	$\{\{ta, ca\}\}$	$\{\{ta, ca\}\}$

Table A4. All Min-Max association rules with $supp_r(*) \geq 0.4$ and $c = 1$.

Numbers	Rule $\phi(supp_r(\phi), c)$
1	$br \wedge to \rightarrow ba$ (0.4653, 1)
2	$ha \rightarrow br$ (0.4257, 1)
3	$to \rightarrow ba$ (0.604, 1)
4	$mi \rightarrow ba \wedge br$ (0.4059, 1)
5	$mi \wedge ba \rightarrow br$ (0.4059, 1)
6	$mi \wedge br \rightarrow ba$ (0.4059, 1)
7	$to \wedge ta \rightarrow ba$ (0.5149, 1)

Table A5. Some Min-Max association rules with $supp_r(*) \geq 0.4$ and $0.8 \leq c < 1$.

Numbers	Rule $\phi(supp_r(\phi), c)$
1	$pr \wedge ba \rightarrow ta$ (0.4059, 0.8723)
2	$pr \wedge ta \rightarrow ba$ (0.4059, 0.9762)
3	$ta \rightarrow br$ (0.604, 0.8133)
4	$pr \rightarrow ba$ (0.4653, 0.8393)
5	$ca \rightarrow ba$ (0.4257, 0.9773)
6	$to \rightarrow ba \wedge ta$ (0.5149, 0.8525)
7	$to \wedge ba \rightarrow ta$ (0.5149, 0.8525)
8	$ba \rightarrow br$ (0.6832, 0.8313)
9	$br \rightarrow ba$ (0.6832, 0.8625)
10	$ba \rightarrow ta$ (0.7327, 0.8916)
11	$ta \rightarrow ba$ (0.7327, 0.9867)
12	$ta \rightarrow ba \wedge br$ (0.5941, 0.8)
13	$ba \wedge br \rightarrow ta$ (0.5941, 0.8696)
14	$ba \wedge ta \rightarrow br$ (0.5941, 0.8108)
15	$br \wedge ta \rightarrow ba$ (0.5941, 0.9836)

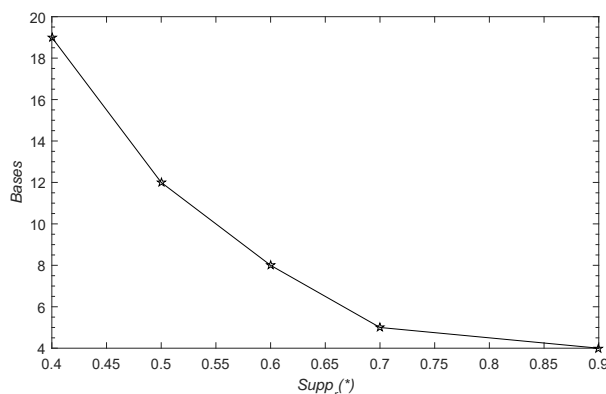


Figure A1. Change of bases in the mushroom database corresponding to $supp_r(*)$.

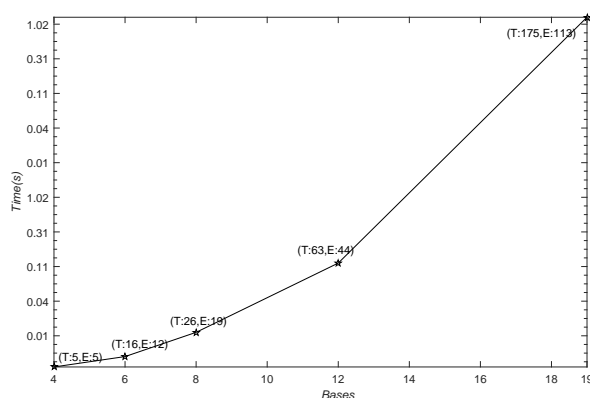


Figure A2. Members of topology, equivalent classes and its generating time in the mushroom database corresponding to bases.

Table A6. All Min-Max association rules in the mushroom database with $supp_r(*) \geq 0.8$ and $c = 1$.

N	Rule $\phi(supp_r(\phi), c)$
1	$gill_attachment_f \rightarrow veil_type_p (0.9742, 1)$
2	$gill_spacing_c \rightarrow veil_type_p (0.8385, 1)$
3	$veil_color_w \rightarrow veil_type_p (0.9754, 1)$
4	$ring_number_o \rightarrow veil_type_p (0.9217, 1)$
5	$gill_attachment_f \wedge veil_color_w \rightarrow veil_type_p (0.9732, 1)$
6	$gill_attachment_f \wedge ring_number_o \rightarrow veil_type_p (0.8981, 1)$
7	$gill_spacing_c \wedge veil_color_w \rightarrow veil_type_p (0.8149, 1)$
8	$gill_attachment_f \wedge gill_spacing_c \rightarrow veil_color_w \wedge veil_type_p (0.8127, 1)$
9	$gill_attachment_f \wedge gill_spacing_c \wedge veil_type_p \rightarrow veil_color_w (0.8127, 1)$
10	$gill_attachment_f \wedge gill_spacing_c \wedge veil_color_w \rightarrow veil_type_p (0.8127, 1)$
11	$veil_color_w \wedge ring_number_o \rightarrow gill_attachment_f \wedge veil_type_p (0.8971, 1)$
12	$gill_attachment_f \wedge veil_color_w \wedge ring_number_o \rightarrow veil_type_p (0.8971, 1)$
13	$veil_type_p \wedge veil_color_w \wedge ring_number_o \rightarrow gill_attachment_f (0.8971, 1)$

Table A7. Min-Max association rules with $supp_r(*) \geq 0.8$ and $0.8 \leq c < 1$.

Numbers	Rule $\phi(supp_r(\phi), c)$
1	$veil_type_p \rightarrow gill_attachment_f$ (0.9742, 0.9742)
2	$veil_type_p \rightarrow gill_spacing_c$ (0.8385, 0.8385)
3	$veil_type_p \rightarrow veil_color_w$ (0.9754, 0.9754)
4	$veil_type_p \rightarrow ring_number_o$ (0.9217, 0.9217)
5	$veil_type_p \rightarrow gill_attachment_f \wedge veil_color_w$ (0.9732, 0.9732)
6	$veil_type_p \rightarrow gill_attachment_f \wedge ring_number_o$ (0.8981, 0.8981)
7	$veil_type_p \rightarrow gill_spacing_c \wedge veil_color_w$ (0.8149, 0.8149)
8	$veil_type_p \rightarrow gill_attachment_f \wedge gill_spacing_c \wedge veil_color_w$ (0.8127, 0.8127)
9	$veil_type_p \rightarrow gill_attachment_f \wedge veil_color_w \wedge ring_number_o$ (0.8971, 0.8971)
10	$gill_attachment_f \rightarrow veil_type_p \wedge veil_color_w$ (0.9732, 0.999)
11	$gill_attachment_f \rightarrow veil_type_p \wedge ring_number_o$ (0.8981, 0.9219)
12	$gill_attachment_f \rightarrow gill_spacing_c \wedge veil_color_w \wedge veil_type_p$ (0.8127, 0.8342)
13	$gill_attachment_f \rightarrow veil_color_w \wedge ring_number_o \wedge veil_type_p$ (0.8971, 0.9209)
14	$gill_spacing_c \rightarrow gill_attachment_f \wedge veil_color_w \wedge veil_type_p$ (0.8127, 0.9692)
15	$gill_spacing_c \rightarrow veil_color_w \wedge veil_type_p$ (0.8149, 0.9718)
16	$veil_color_w \rightarrow gill_spacing_c \wedge veil_type_p$ (0.8149, 0.8354)
17	$veil_color_w \rightarrow gill_attachment_f \wedge gill_spacing_c \wedge veil_type_p$ (0.8127, 0.8332)
18	$veil_color_w \rightarrow gill_attachment_f \wedge veil_type_p$ (0.9732, 0.9977)
19	$veil_color_w \rightarrow gill_attachment_f \wedge ring_number_o \wedge veil_type_p$ (0.8971, 0.9197)
20	$ring_number_o \rightarrow gill_attachment_f \wedge veil_type_p$ (0.8981, 0.9744)
21	$ring_number_o \rightarrow gill_attachment_f \wedge veil_color_w \wedge veil_type_p$ (0.8971, 0.9733)
22	$gill_attachment_f \wedge veil_color_w \rightarrow ring_number_o \wedge veil_type_p$ (0.8971, 0.9218)
23	$gill_attachment_f \wedge veil_color_w \rightarrow gill_spacing_c \wedge veil_type_p$ (0.8127, 0.8351)
24	$gill_attachment_f \wedge ring_number_o \rightarrow veil_color_w \wedge veil_type_p$ (0.8971, 0.9989)
25	$gill_spacing_c \wedge veil_color_w \rightarrow gill_attachment_f \wedge veil_type_p$ (0.8127, 0.9973)
26	$gill_attachment_f \wedge veil_type_p \rightarrow veil_color_w \wedge ring_number_o$ (0.8971, 0.9209)
27	$gill_attachment_f \wedge veil_type_p \rightarrow ring_number_o$ (0.8981, 0.9219)
28	$veil_type_p \wedge veil_color_w \rightarrow gill_attachment_f \wedge ring_number_o$ (0.8971, 0.9197)
29	$veil_type_p \wedge ring_number_o \rightarrow gill_attachment_f \wedge veil_color_w$ (0.8971, 0.9733)
30	$veil_type_p \wedge ring_number_o \rightarrow gill_attachment_f$ (0.8981, 0.9744)
31	$gill_attachment_f \wedge veil_type_p \wedge veil_color_w \rightarrow ring_number_o$ (0.8971, 0.9218)
32	$gill_attachment_f \wedge veil_type_p \wedge ring_number_o \rightarrow veil_color_w$ (0.8971, 0.9989)
33	$gill_attachment_f \wedge veil_type_p \rightarrow gill_spacing_c \wedge veil_color_w$ (0.8127, 0.8342)
34	$gill_spacing_c \wedge veil_type_p \rightarrow gill_attachment_f \wedge veil_color_w$ (0.8127, 0.9692)
35	$veil_type_p \wedge veil_color_w \rightarrow gill_attachment_f \wedge gill_spacing_c$ (0.8127, 0.8332)
36	$gill_attachment_f \wedge veil_type_p \wedge veil_color_w \rightarrow gill_spacing_c$ (0.8127, 0.8351)
37	$gill_spacing_c \wedge veil_type_p \wedge veil_color_w \rightarrow gill_attachment_f$ (0.8127, 0.9973)
38	$gill_attachment_f \wedge veil_type_p \rightarrow veil_color_w$ (0.9732, 0.999)
39	$veil_type_p \wedge veil_color_w \rightarrow gill_attachment_f$ (0.9732, 0.9977)
40	$gill_spacing_c \wedge veil_type_p \rightarrow veil_color_w$ (0.8149, 0.9718)
41	$veil_type_p \wedge veil_color_w \rightarrow gill_spacing_c$ (0.8149, 0.8354)

Table A8. The number of approximate rules and reduction in the mushroom database with confidence $c \geq 0.5$.

$supp_r(*)$	TRX	Min-Max (N, Rt)	Reliable (N, Rt)	TR	N-Min-Max (N, Rt)
0.4	2528	(465, 0.82)	(361, 0.86)	1825	(420, 0.77)
0.5	835	(175, 0.79)	(135, 0.84)	514	(190, 0.63)
0.6	228	(59, 0.74)	(52, 0.77)	136	(65, 0.52)
0.7	161	(39, 0.74)	(34, 0.79)	90	(41, 0.54)
Average		0.77	0.82		0.62

Table A9. Number of approximate rules and reduction in the connect-4 database with confidence $c \geq 0.5$.

$supp_r(*)$	TRX	Min-Max (N, Rt)	Reliable (N, Rt)	TR	N-Min-Max (N, Rt)
0.94	199,560	(49,407, 0.75)	(10,220, 0.95)	88,116	(7914, 0.91)
0.95	77,206	(24,794, 0.68)	(5245, 0.93)	39,768	(4731, 0.88)
0.96	26,856	(11,452, 0.57)	(2538, 0.91)	16,356	(2535, 0.85)
0.97	7895	(4439, 0.44)	(1214, 0.85)	5690	(1294, 0.77)
Average		0.61	0.91		0.85

Table A10. Templates, association rules and generalized association rules in the connect-4 database.

Template 1	$A = \{d_5 - b, d_6 - b, e_5 - b, f_5 - b, f_6 - b\} (supp_r(A) = 0.9573)$
ψ_1	$d_5 - b \wedge d_6 - b \wedge f_5 - b \wedge f_6 - b \rightarrow e_5 - b (0.9573, 0.99)$
Reduction of ψ_1	$d_5 - b \wedge f_5 - b \rightarrow e_5 - b (0.9573, 0.99)$
Rule 1	$\phi \equiv a_6 - b \wedge e_4 - b \wedge f_5 - b \rightarrow e_6 - b \wedge f_6 - b (0.9502, 0.9694)$
	$\phi_{GA} \equiv a_6 - b \rightarrow e_6 - b \wedge f_6 - b (0.9502, 0.9576)$
	$\phi_{ga} \equiv a_6 - b \wedge e_4 - b \wedge e_5 - b \wedge f_5 - b \rightarrow e_6 - b \wedge f_6 - b (0.9502, 1)$

Table A11. The number of approximate rules and reduction in the chess database with confidence $c \geq 0.5$.

$supp_r(*)$	TRX	Min-Max (N, Rt)	Reliable (N, Rt)	TR	N-Min-Max (N, Rt)
0.90	10,614	(8371, 0.21)	(2483, 0.77)	9230	(2034, 0.78)
0.91	5785	(5050, 0.13)	(1571, 0.73)	5354	(1357, 0.75)
0.93	2338	(1948, 0.17)	(688, 0.71)	2110	(648, 0.69)
0.95	468	(459, 0.02)	(196, 0.58)	466	(195, 0.58)
Average		0.13	0.70		0.70

Table A12. Templates, association rules and generalized association rules in the chess database.

Template 1	$A = \{mulch_f, skach_f, wkna8_f\} (supp_r(A) = 0.9002)$
ψ_1	$mulch_f \wedge skach_f \rightarrow wkna8_f (0.9002, 0.9492)$
Rule 1	$\phi \equiv qxmsq_f \wedge spcop_f \wedge wkna8_f \rightarrow bkon8_f \wedge hdchk_f \wedge reskd_f \wedge thrsk_f (0.8326, 0.9103)$
	$\phi_{GA} \equiv qxmsq_f \wedge spcop_f \rightarrow bkon8_f \wedge hdchk_f \wedge reskd_f \wedge thrsk_f (0.8827, 0.9107)$
	$\phi_{GC} \equiv qxmsq_f \wedge spcop_f \wedge wkna8_f \rightarrow bkon8_f \wedge hdchk_f \wedge reskd_f (0.8698, 0.9136)$
	$\phi_{GAC} \equiv qxmsq_f \wedge spcop_f \rightarrow bkon8_f \wedge hdchk_f \wedge reskd_f (0.9215, 0.9507)$
	$\phi_{ga} \equiv qxmsq_f \wedge spcop_f \wedge wkna8_f \wedge stlmt_f \rightarrow bkon8_f \wedge hdchk_f \wedge reskd_f \wedge thrsk_f (0.8326, 0.9103)$

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